INDUSTRIAL INFORMATICS

and EMBEDDED SYSTEMS

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COURSE SECTIONS:

- DIGITAL INTERFACING
- INDUSTRIAL COMMUNICATIONS & FIELD BUSES
- EMBEDDED ARM MICROPROCESSORS



THE EXAM

Typically: 1) a co

1) a code fragment in ARM assembly language to be interpreted

2) an exercise on the evaluation of the frequency of the pulses provided by a suitable sensor, considering a microprocessor whose bits are employed for different purposes (driving a motor, acquisition of a physical a magnitude, generating a signal, PWM, PFM, ...)

As the case is:1) questions on the different parts of the course to establish if mechanisms
have been really captured
2) exercises on the communication protocol part



Companies that propose thesis:

MARELLI MOTORSPORTS (Corbetta Milano): video data logging, telemetry

TEMIS (Corbetta Milano): automotive, satellites, data acquisition and communication

BDSOUND (Assago): sound engineering, audio signals acquisition through microcontrollers (Cortex M4, ST, ...)

AZCOM (Rozzano): wireless (and not) safe communications through sw control (DSP and FPGA processors)

FACILITY LIFE: cybersercurity

ST MICROELECTRONICS: embedded processors

DSP, GPU, FPGA technologies applied to:

Brain cerebellar simulation

Hyperspectral image processing for high accuracy cancer detection

Glaucoma identification in ophtalmic imaging

Automotive applications

Signal processing through embedded microprocessors



FOCUS OF THE COURSE



- •**Control algorithm** = the controller operating principle
- •**Process** = the problem. Know the transfer function
- •**Minimize** the gap between SP and PV (and dependency on environmental variables)
- •Oscillations due to the velocity of the answer to SP variations \Rightarrow minimize the transient
- **Overshootings** \Rightarrow limit the amplitude
- •Stability (no infinite or diverging oscillations)

•Delays

Gain



- INDUSTRIAL COMMUNICATIONS & FIELD BUSES

- EMBEDDED ARM MICROPROCESSORS



Interfacing = interactions between active operative devices (i. e. CPU) and/or passive ones (i. e. memories, sensors, ports)

The course is mainly (but not only) focused on the interactions between **calculator** and **external world** (every thing not directly manageable by the CPU through buses)

Several kinds of **in-homogeneities**:

- *Physical*: the external world is characterized by physical parameters that can be represented only through electric entities with suitable amplitudes (transducers)
- Formal: codification (BCD/Gray), representation (pulse count or frequency measurement), signal processing/management (PWM or PFM modulation)
- Temporal: how external phenomena are synchronized with signal acquisition/emission? (timers, interrupt/polling/DMA, real time protocols)
- Spatial: computers and controlled devices are not physically adjacent (transmission)



• *Attitudinal*: Human Computer Interaction

A suitable **chain** of interacting devices allows to overtake these in-homogeneities

Typical computer-external world interfacing

Interface design = identification of a chain of components able to wholly satisfy the requirements of data transformation and communication





A general interface digital circuit

- •I/O buffer (Fifo)
- *State registers* for significant situations (errors, overrun, interrupt ...)
- *Command registers* to setup mode functions and modify state bits
- Various functions: DMA, ser/par transfers

ADDRESSING ON THE BASIS OF A MEMORY MAP



ADDRESSING ON THE BASIS OF A MEMORY MAP



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On off signals

- •They represent physical entities acquired through single bits of a microprocessor port
- •What is important is the value of the bit, not its variation
- •8-16-32 bit groups acquisition (bytes, words)
- •Let's **test** the Kth bit



•And now two assembly routines (the right one is optimized maybe by the compiler)



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On off signals

•Let's **set** the Kth bit

```
Esempio in C

if (var_bool)

immagine |= mask_K; // OR con mask

else

immagine &= ~ mask_K; // AND con COMPLEMENTO di mask
```

Esempio in Assembler					
MOV	AL, MASK_K	; maschera			
AND /	AL, [VAR_BOOL]	; isola bit in posizione K			
MOV	AH, AL	; salva in AH			
MOV /	AL, MASK_K				
CPL	AL	; maschera complementata			
AND /	AL, [IMMAGINE]	; azzera bit in posizione K			
OR	AL, AH	; inserisce nuovo bit in posizione K			
MOV	[IMMAGINE], AL	; aggiorna variabile immagine			

On off signals: "state" acquisition

• The value of a bit is read to establish the level (*state*) of a line related to a physical entity: use of a *three state buffer*



On off signals

• three state buffers to de-couple the line with respect other devices



COTS buffer components (Texas Instr. SN 74AC244)

10E	1	U 20	Vcc
1A1	2	19	20E
2Y4	3	18	1Y1
1A2	4	17	2A4
2Y3	5	16	1Y2
1A3	6	15	2A3
2Y2	7	14	1Y3
1A4	8	13	2A2
2Y1	9	12	1Y4
GND	10	11	2A1

· ,











2OE

2A3



On off signals: event triggering

•The value of a bit is read just to establish if a commutation has occurred (event) of a line related to a physical situation, by comparing it with the previous value

•2 possible approaches:

-SW triggering through polling

>A microprocessor port is continuously sampled. The software compares the new value with the old one and identifies the eventual commutation

-triggering through interrupt

>The signal carrying the "event" is connected to a microprocessor port's pin

- >The single commutation edge is determined through pins edge triggering (falling/raising edges)
- >Both the edges are acquired through suitable electronic circuits sensible to edge or to levels. Where is the difference?



On off signals: "event" acquisition through interrupt

Latch (sensible to the level)

FF (sensible to the edge)



On off signals

• a bit value is read to establish the possible commutation using *one-edge sensible circuits*. Below a circuit able to transform commutation of **both** the edges in a single direction



On off signals

a bit value is read to establish the possible commutation using *level sensible circuits* (a FFJK, or a SR, is required to keep the input fixed)



Physical signal adjustment: opto-couplers

- •Mandatory if grounds are different or if possible over-voltages
- $\bullet R_L$ to determine necessary current for LED power up

•R_P pull-up (saturation vs "steep" edges: what possible values?)



DIGITAL INTERFACING: FILTERING

On off signals

- •This kind of signals are produced by relays, electronic switches, buttons, limit-switches
- •Possible *spurious* states irrelevant if the signal level is required but misleading if we are interested in the commutation (i. e. an associated *event*)



- *electric* vs *algorithmic* filtering
 - *–electric*: required if interrupt o counters are used (\Rightarrow RC + Schmitt trigger)

–algorithmic: necessary if the signal is sampled. The sampling will turn out in a series of values among which some are significant others not (spurious). The choice among them is driven by the information that the signal carries \Rightarrow definition of a suitable *transient death time*

Fine corsa a leva, stelo, pulsante Lever, stem, push-button,



Fine corsa ad albero

tree limit switches



Anello di tenuta ermetica a tre vie per l'albero in fluoroelastomero brevettato da Allen-Bradley rinforzato per sopportare lavaggio con getto in pressione fino a 83 bar (8.270 kPa)

La guarnizione a diaframma a tenuta stagna fra la testa operatrice e il corpo del fine corsa offre ulteriore protezione contro agenti contaminanti quanto la testa operatrice viene rimossa

In grado di sopportare immersione in condizioni che superano i requisiti NEMA 6P

o con attacco da 30 mm x 60 mm con corpo metallico dotati di albero e componenti in acciaio inossidabile oltre a nuovi materiali per una migliore resistenza alle sostanze chimiche.

1 modelli 802MC sono

Electric filter through RC circuit + Schmitt trigger

Choose carefully the values of R and C, since they influence the circuit time constant and thus the reliability of the microprocessor acquisition

Slow capacitor charge/discharge could imply a slow crossing in the uncertainty range, bringing an oscillating behaviour. Therefore a Schmitt trigger could be suitable.





Algorithmic filtering: first edge detection

- •The first variation is considered as really significant and accepted
- •Further variations are ignored during the time interval T (*transition death or dead time*)
- •The algorithm does not work with isolated spurious signals
- •Spurious signals follow a real significant event and run out within the T interval. Ok in case of mechanical contacts



First edge detection, a simple assembly code

; Routine da chiamare ciclicamente ad ogni periodo di campionamento. : Riporta in AL il valore filtrato ACQ BIT N: CMP [TRANSITORIO].0 JE REGIME ;qui in transitorio - conta il tempo trascorso DEC [TRANSITORIO] JMP ESCI ;qui transitorio esaurito - pronto a sentire event. commutaz. REGIME: IN AL.PORTA IN ;legge la porta AND AL.MASK N isola il bit considerato: CMP AL,[PRECEDENTE] :confronta con valore preced. JE ESCI ;qui rilevata commutazione MOV [PRECEDENTE],AL aggiorna il valore [TRANSITORIO],T_MORTO rilancia il transitorio: MOV ESCI: MOV AL.[PRECEDENTE] riporta il valore valido; RET

Algorithmic filtering: expired transient detection

- •The first variation is not significant
- •The corresponding state is accepted as valid only if unchanged within T time interval
- •The algorithm works well with isolated spurious signals
- •A low pass filtering effect (delay)



Exhausted transient detection, a simple assembly code



On off signals emission

- •The output on a μ P port is carried out through latches (permanent values)
- •Output of initial setup values or after a reset (0 if possible)
- •Or simply those values that will manipulate the process through the actuator
- •Output through "image" variables (it is not possible to output single bits)
- •Image variables processed through AND/OR/XOR masks
- the National Instruments Latch 54ACT374 component



Logic Diagram



Please note that this diagram is provided only for the understanding of logic operations and should not be used to estimate propagation delays.



Pin Assignment for LCC

Truth Table

Inputs			Outputs
Dn	CP	OE	On
Н	ł	L	Н
L	~	L	L
х	х	Н	Z

H = HIGH Voltage Level L = LOW Voltage Level X = Immaterial Z = High Impedance

IOW-to-HIGH Transition

Latch 54ACT374

National Instruments



Pulses can be acquired for:

- Temporization
- •Counting a suitable number of moving objects
- Calculating object position
- Calculating object velocity

Binary digital signal commutations to which a single event is associated considered as significant

Sometimes both the edges are useful since associated to the beginning and the end of a event \Rightarrow interval measurement

Pulses can be emitted for:

- •PFM device driving
- •PWM device driving
- •Power supply to an electric load through partialization of the signal

Pulse acquisition for counting (object passage, complete rounds ...)





- •LED + phototransistor to reveal the object passage
- •Total primary counter must be upgraded (i. e. program variable)
- •Max reachable value => type definition => unsigned long (32 bit)
- •Non volatile counting (eventually secondary counters in powered buffer memories / files. Trade-off between freshness and CPU overload)
- Manual vs automatic reset operations

•Variables must be stored through atomic saving operations in memory

Hw primary counter

- •Less sw elaboration charge onto the CPU
- •Mandatory if pulse frequency not acquirable through polling from the μ P port (it requires few assembly instructions)
- If the counter is connected with a μP port we read a state information
- •If the counter is connected with an interrupt line (triggers the interrupt) we read an event information



Hw primary counter: a commercial component Texas SN74193





•Clear pin to reset

- •Up/Down pins for setting bidirectional inputs
- Load pin to values preset
- •A, B, C, D preset inputs
- •Borrow/carry pins

External (primary) hw counter providing a "status" information

- N bit hw counter
- µprocessor port connection
- •Acquisition routine periodic activation T_c
- where $T_c \leq (2^N-1)*T_{imp}$
- Input from the port (present counter AGGIORNA_CONT. value) IN AL, P
 - # pulses between two inputs (V_{new} V_{old}
 - added to a total (secondary) counter (carry management)
- AGGIORNA_CONTATORE: SUB AX,AX IN AL,PORTA_CONT SUB AL,[PRECEDENTE] ADD [PRECEDENTE],AL ADD [CONTATORE],AX ADC [CONTATORE + 2],0 RET
External (primary) hw counter providing a "status" information

- If n° of bits > than readable by a single instruction
- •Bytes from the counter are read with successive read phases
- Auxiliary register with the same width of the counter to be read in two phases



AGGIORNA CONTATORE: IN AL, P MSBYTE RIPETI: MOV BH.AL IN AL.P LSBYTE MOV BL.AL IN AL.P MSBYTE CMP AL.BH JNE RIPETI ;qui BX = valore valido MOV AX.BX AX,[PRECEDENTE] SUB [PRECEDENTE].AX ADD ADD [CONTATORE],AX ADC [CONTATORE + 2],0

External (primary) hw counter providing an "event" information

Interrupt line connected to a single counter bit or to the carry bit
Used as "pre-scaler" (frequency division)

 $F_{out} = F_{in} / 2^{N} (if carry bit connected)$ $F_{out} = F_{in} / 2^{k} (if k^{th} bit connected)$

Memory SW counter upgrading (if high frequencies)
Least significant count directly read from the counter

SW counter Direc

Direct acquisition

Pulses to evaluate an object position

•Pulses corresponds to incremental movements. Where is the "zero" point?

- -Is it an un-surmountable extreme? Unsigned numbers
- -Is it an arbitrary position? Signed numbers
- •The incremental encoder



Incremental Encoder

- •Two concentric crowns, 1/4 period out of phase
- •Two couples of LED/transistor (1 light, 0 dark).
- Minimum measurable angle = half window





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Counting specs

- definition of "zero" (natural/integer counting variable)
- maximum counting vs granularity (size)

Counting

- edge detection
- movement direction
- primary counter upgrading (then eventually upgrade the secondary one)
- 2 possible approaches: sw/hw



A	В	C
\mathbf{P}	0	+1
\mathbf{P}	1	-1
N	Ο	-1
N	1	+1
Ο	Ρ	-1
1	\mathbf{P}	+1
Ο	Z	+1
1	Ζ	-1

- •A and B: 2 pulse trains connected to two pins of the μP port
- Commutations detection
- Movement direction identification: forward (+) or backward (-)
- Counter variables (primary and eventually secondary)
- algorithm (hypothesis clockwise rotation considered as positive)



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An example of sw management of a bi-directional incremental encoder

AGGIORNA POS ENCODER: AL.P ENC IN AND AL.3 :maschera i bit utili MOV BL.AL SHR BL.1 :BL.0 = B (porta segnale B in bit 0 XCHG AL, [PRECEDENTE] XOR AL.[PRECEDENTE] JΖ ESCI :qui rilevata una commutazione AND AL.1 MOV CL.AL :CL.0 = 1 se commutato A MOV AL,[PRECEDENTE] AND AL.1 :AL.0 = A AL.BL XOR AL.CL XOR DECREMENTA JNZ ;qui si deve incrementare ADD [CONTATORE],1 JMP ESCI DECREMENTA: SUB [CONTATORE],1 ESCI: AL = A 🕀 CL = Anew AL CL ACTION rotation notes RET B B个 Anew=B CW 0 0 0 +1 B↑ Anew≠B CCW 1 0 1 -1 0 1 1 -1 A个 Anew=B CCW University of Pa 1 0 A↑ Anew≠B CW +1

Routine activation frequency?



We have to distinguish between: case ok B(t-1)≠B(t) with A=1, (cw) case ko B(t-1)≠B(t) with A=0 (ccw) F_{clk}= 2*F_s This on both the A e B signals Fclk= 4*Fs

Hw management of a bi-directional incremental encoder



Activation frequency?

The hw logic must distinguish among commutations of the same signal (A o B); however, within the present commutation and the next one it must also recognise if the other signal is 0 or 1.

$$F_{clk} = 8 * F_s$$

Hw management of a bi-directional incremental encoder



Hw management of a bi-directional incremental encoder



Hw management of a bi-directional incremental encoder



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Hw management of a bi-directional incremental encoder

CCW rotation



Pulses for timing

Time-stamping means to associate events and times into which they happen

Errors due to the model describing time evolution within the machine (TIG – temporal granularity)

Delays due to timer acquisition sw (T_{el}) and to event recognition (T_{rev})

Latencies = time interval between the event and its perception (T_{lat}) . Non deterministic, if they are due to resources charged with task perception and in that moment devoted to other activities (priority). Typical in interrupt service and in processes activation



Pulses for timing (maximum priority routine)

Events internal to the machine (i. e. switch of a control unit signal)

 $t-TIG+T_{el} < T_{st} < t+T_{el}$

Events external to the machine (i. e. switch of a signal read through a port)

 $t-TIG+T_{rev}+T_{el} < T_{st} < t+T_{lat.max}+T_{el}+T_{rev}$

Measurement of an interval between external events (constants components neutralise each other, deterministic ones do not disappear)

$$-TIG-T_{lat.max} < (T_{st1}-T_{st2}) < TIG+T_{lat.max}$$

What is TIG? It depends on the machine time representation model. Let us open a parenthesis ...

- Is the time an event container (box) or the events cause the course of the time?
- During life we follow the first approach, in computer science we are close to the second one
- CLOCKS = phenomena that generate regular events (ticks) used to quantify the time passing
- We need a model

TIME AS IT IS REPRESENTED WITHIN A CPU

- •Hp: bijective correspondence among time moments and real numbers
- •Continuity, limit, derivative, integral existence ecc.
- •Time as linear, monotonic, discrete
- •Clock = time events generator with domain N⁺
- •Quantization = duration of the phenomenon taken as clock, usually very negligible with respect events that must be measured

TIMING MODEL

•LOCAL TIME UNIT (UTL) = absolute time interval corresponding to 1 in the measurement unit of the clock

•LOCAL TIME BASE (BTL) = multiple of the UTL, representing the period of the cyclic phenomenon assumed as clock

•Let' suppose that in a PC the real time clock is 55 msec: the time is represented in msec (UTL), but it is increased every 55 (BTL)

```
•UTL = resolution, BTL = granularity

•If BTL=1, resolution = granularity

•If T = clock time and t = absolute time,

T=n*BTL

n*BTL*UTL \le t \le (n+1)*BTL*UTL
```

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TIME MODEL: PRECISION ERRORS

•They are due to how the phenomenon (assumed as clock) differs from the ideal behavior

•PRECISION

A clock is precise if t(tick(n)) - t(tick(n-1)) = BTL*utl + e(n) with e(n)=0

– Systematic error \Rightarrow the phenomenon taken as clock features a not correct frequency

- Null mean errors \Rightarrow not stable frequency (jitter)
- precision measurement = BTL*utl /mean_err = (10⁴-10⁶)

•CORRECTNESS

A clock is correct if abs(T-t/UTL) < BTL, that is the local time T is synchronised with the absolute time (within the quantization error BTL)

-A precise clock if is correct in t, is correct in t'>t

-Sistematic unaccuracy (mean_err≠0) ⇒ no correctness

-If intervals are measured, systematic errors do not exert any influence

TIME MODEL: QUANTIZATION ERRORS

•They are caused by the representation of the time

- •ORDER: it is not possible to establish an order between 2 events that happen within 2 consecutive ticks (i. e. within the same BTL)
- •CONTEMPORANEITY: if the time interval between two events is $\Delta T < BTL$, the computer considers them as simultaneous (contemporaneous)

•MEASUREMENT:

- 2 events, whose real time distance is ΔT , are considered as within N*BTL < M(ΔT) < (N+1)*BTL N=floor(ΔT /BTL)
- events, whose time distance is N*BTL, can correspond to a real time interval ∆T: (N-1)*BTL*UTL < ∆T < (N+1)*BTL*UTL



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TYPICAL TIMING SYSTEMS (TIMERS)



- Prescaler (frequency division, $2^{k}=1-256$)
- TIMER = N bit binary cyclic counter (16-24)

TIME BASE COUNTER = 1 tick when the set period (TBV) ends + restoration TBV Decrement counters

TYPICAL TIMING SYSTEMS (TIMERS)

• <u>TIMER section</u>

- TIMER RESOLUTION (count units) = $2^{k}/F=UTL$
- TIMER PERIOD = $2^{N} * 2^{k}/F$
- BTL= 2^{N} (on the overflow)



(COUNT+TBV) mod2^M=COUNT+TBV-2^M

 $= 2^{M} - x + TBV - 2^{M} = TBV - x$

to include counts lost due to interrupt latency

- TIME BASE section
- TIME BASE COUNTER RESOLUTION = $2^{k}/F=UTL$

TIME BASE COUNTER PERIOD = TBV $* 2^{k}/F$

BTL=TBV

TBV RESTORE

- •Set TBV automatically or ...
- •Explicit setup within the sw routine that manages tick interrupt but ...
- ●(CONTEGGIO+TBV) mod2^M

SOME IDEAS

F=1 MHz, Prescaler=1 \Rightarrow period 1µsec, if N=16 \Rightarrow 65000 µsec

F=1 MHz, Prescaler=256 \Rightarrow period 256µsec, if N=16 \Rightarrow 16 sec

F=20 MHz, Prescaler=256 \Rightarrow period 12.5 μ sec, if N=24 \Rightarrow 200 sec

Process activation	\Rightarrow TIME BASE COUNTER (TBC)
Time stamping (date and time)	\Rightarrow tick TBC
	\Rightarrow overflow TIMER
Interval measurement	\Rightarrow TIMER

NB. If long intervals are measured we need a sw variable onto which to store the counter values. This implies a partial un alignment between the total sw counter value and the hw register (containing something like a «fractional part» of the overall value).

Pulses are acquired to calculate the velocity (i. e. frequency) of an object

Angular velocity or (scale factor) linear velocity can be determined

Problem data

- Encoder granularity = (round or meter/pulse)⁻¹ => quantization effect
- Maximum and minimum frequency to be measured
- Pulse (usually given by a sensor) duration

Specs

- Measurement production frequency (acceptable depending on the particular phenomenon to be measured)
- Delay between the acquisition instant and the moment to which it is related
- Errors (relative vs. absolute, related to measurement value or to the full scale value, Erm% vs. Erf%)

VELOCITY CALCULATION

Pulses are counted during a known time interval





What is measured is an "average velocity"

Pulse acquisition techniques

Software

Cyclic (Tc period) pulse acquisition If time for reading negligible and Tg granularity of the used clock (BTL* UTL) Error on the event detection: -Tg<E<Tc

Interrupt

The signal carrying the pulse is connected to the interrupt pin of a microprocessor T_{lat.max} = max. latency of the interrupt service routine Te = execution time of the interrupt service routine Error on the event detection: -Tg+Te<E<T_{lat.max}+Te

Hardware

Dedicated circuits with working times negligible with respect to event duration

Error on the event detection: -Tg<E<0

Pulse count during time interval: I-SW, T-SW

•Sw routine executed with period Tc:

- Input port read
- if positive (raising up) edge I count++;
- if positive edge T vel=(enc. granularity)*counter_val/ Δ T; count=0;

Constraints

Tc< minimum pulse duration \Rightarrow f Tc< 1/2 (f = pulse frequency) f_P (measure production freq.) = 1/ Δ T; f_{min} (minimum detectable frequency) > 1/ Δ T

•Errors

DTM= real (effective) count interval; n_p =real n° of pulses DTM f = real signal frequency = n_p /DTM; f_m = measured frequency = n (counter value) / Δ T

Quantization err. on pulse identification [-1 ... 1] \Rightarrow n_p -1 < n < n_p +1 Time err. on ΔT estimation [-Tc ... +Tc] $\Rightarrow \Delta T$ -Tc < DTM < ΔT + Tc

$$f(\Delta T - Tc) - 1/\Delta T < f_m < f(\Delta T + Tc) + 1/\Delta T$$

$$Ea_{sup} = (1+fTc)/\Delta T = -Ea_{inf}$$

 $Er_{sup} = Ea/_{fmin} = -Er_{inf}$

Ρ

Pulse count during time interval: I-HW, T-int

•N bit counter connected to a microprocessor port, read during interrupt routine at the beginning and at the end of ΔT . Steps of the interrupt service routine:

- interrupt disable and context saving
- count = present counter value previous value;
- freq = count/ ΔT ;
- previous value = present counter value;
- context restore and interrupt enable

Constraints

$$f_{max} \Delta T < 2^{N} (f=frequency of pulses); f_{P} = 1/\Delta T; f_{min} > 1/\Delta T$$

•Errors

Quantization error on pulses identification [-1 ... 1] Time error on ΔT estimation [-T_{Imax} ... +T_{Imax}]

$$f(\Delta T - T_{lmax}) - 1/\Delta T < f_m < f(\Delta T + T_{lmax}) + 1/\Delta T$$

$$Ea_{sup} = (1+fT_{Imax})/\Delta T = -Ea_{inf}$$



 T_{Imax} = max. interrupt latency

$$Er_{sup} = Ea/_{fmin} = -Er_{inf}$$

Pulse count during time interval: I-HW, T-HW



Constraints

 $f_{max} \Delta T < 2^{N}$ (f= pulse frequency); $f_{P} = 1/(\Delta T + T_{lmax})$; $f_{min} > 1/\Delta T$

•Errors

Quantization error on pulses identification [-1 ... 1] Time error on ΔT estimation [-T_{ck} ... 0]

$$f(\Delta T - T_{ck}) - 1/\Delta T < f_m < f\Delta T + 1/\Delta T$$

Considerations on pulse count during time interval

- •Typical mode: interrupt for time, pulse count with counter
- for a good resolution $\Delta T >>$ pulse period
- •Constant quantization error on frequency measurement $(1/\Delta T)$

 $\bullet \Delta T$ has to be chosen considering the minimum frequency to be measured and the relative error

TRADEOFF ON ERRORS

- • $\uparrow \Delta T$, \Downarrow errors and measurable frequencies \Downarrow , \uparrow measureent delay
- $\bullet \Downarrow T_{ck}, \Downarrow \text{ errors,} \Uparrow \text{ timer bits}$
- $\bullet \Downarrow \mathsf{T}_{\mathsf{latency\,int.}} \Downarrow \mathsf{errors,} \Uparrow \mathsf{process}$ management complexity
- • \Downarrow Tc \Downarrow errors and measurable frequencies \Uparrow , \Uparrow CPU load

Measure of the interval between two consecutive pulses



 $T_{ck} 2^N > 1/f_{min}$; $f_P = f$ to measure (minimum f_{min}); $f_{max} < 1/T_{ck}$

•Errors

 $f/1+f(T_{ck}+T_{lmax}) < f_m < f/1-f(T_{ck}+T_{lmax})$

Considerations on the measurement of the interval between two pulses

- typical mode: that showed
- null frequency measurement is critical
- variable acquisition period (f_p) and measurement production
- minimum resolution when frequency is high, max. when it is low
- measur. interval = N pulses, with N set depending on the chosen resolution
- ΔT chosen on the basis of the minimum frequency

COMPROMESSI SU ERRORI

 $\downarrow \mathbf{T}_{ckr} \downarrow \text{ errors, } \uparrow \mathbf{f}_{max} \text{ detectable, } \uparrow \text{ timer bits}$ $\downarrow \mathbf{T}_{latency int.} \downarrow \text{ errors, } \uparrow \mathbf{f}_{max} \text{ detectable, } \uparrow \text{ complexity}$

- The frequency of a pulse train must be measured in the variable range of [5 KHz ... 500 KHz]. The measurement must be provided every ΔT sec, allowing a small tolerance in the periodicity due to the real availability of the CPU:
- •Describe how the chosen technique works and choose suitable devices so as the uncertainty range affecting the measure in the worst case is [-0.15% ... + 0.1%];
- •Verify that the constraints required by the chosen technique are satisfied;
- •Determine the maximum and minimum frequencies really measurable with the imposed constraints;
- •Calculate the uncertainty range in the frequency measurement if the quantization error on the ΔT interval estimation is supposed to be equal to
- +/- "quantum".

Pulse acquisition for synchronising

events corresponding to a valid information



• peripheral ready to receive



• Data_valid; in A, (PORTA_PIO)



• Strobe from peripheral; out A, (PORTA_PIO)



PWM PULSE EMISSION

- Control of processes with a slow behavior that are sensible to the average value of the control signal applied (or of the manipulated variable). For example the environmental temperature
- A pulse with A amplitude and D(t) duration where 0<D(t)<P (period)
- Average value $(V_m) = A^*D/P$ or $A^*DC(t)$ DC = duty cycle (D/P)
- Relative quantization error on V_m is P_{ck} (clock granularity) / P



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EXAM TEXT (July 2006)

- Design a system able to heat up an incubator at 20 °C for 2 minutes, using a typical microprocessor equipped with a parallel PIO port connected to a buffer and to a amplification stage with gain A=10
- Suppose that the ratio between environmental temperature and the average applied power (thermal efficiency) is 80 °C/W and that the heater can be electrically modeled as a 1 KΩ resistance.
- Suppose to heat for 10 minutes at 30°, for 5 minutes at 20° and for 5 minutes at 10°: what is the final temperature reached?


EXAM TEXT (July 2006)



$$W_{EL} = \frac{V^2}{R} = 0,625 W \qquad Temp = 0,625 \cdot 80 = 50 \ ^\circ \mathbf{C}$$
$$W_{EL_{20^\circ C}} = \frac{20}{80} = 0.25W \qquad V_{20^\circ C} = \sqrt{1000 * 0,25} = \sqrt{250} = 15,8$$
$$\frac{\Delta T}{120} = \frac{15,8}{25} \rightarrow \Delta T = 76,8s$$

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$$\frac{T1 \cdot \Delta T_1 + T2 \cdot \Delta T_2 + T3 \cdot \Delta T_3}{\Delta T_1 + \Delta T_2 + \Delta T_3} = \frac{30 \cdot 10 + 20 \cdot 5 + 10 \cdot 5}{20} = 22,5^{\circ}C$$

$$T_1 = 30 \rightarrow W_{EL1} = \frac{30}{80} \rightarrow V_1 = \sqrt{1000 \cdot \frac{30}{80}} = 19,36 V$$

$$T_{2} = 20 \rightarrow W_{EL1} = \frac{20}{80} \rightarrow V_{2} = 15,8 V$$

$$T_{3} = 10 \rightarrow W_{EL1} = \frac{10}{80} = 0,125 \rightarrow V_{1}$$

$$D_{1,2,3} = \frac{V_{1,2,3}}{V_{max}} \Delta t_{1,2,3}$$

$$= \sqrt{125} = 11,2 V$$

 \sim

- Transformation of electric energy in mechanic (rotation vel. and torque)
- If the voltage supply is fixed a dc motor absorbs a current proportional to the exerted torque (thus it depends on the applied load)
- When the motor is in *stale* it absorbs much more current than in the usual case so to avoid damages is important to define two "limit" magnitudes
- **Stalemate current**: max current absorbed by a motor at the nominal voltage, during the stalemate phase (from hundreds of mA to few A).
- **Stalemate torque**: the torque provided when the driveshaft is blocked (at the startup or due to an external load), with a nominal voltage and maximum current (100 Nm)



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- Sometimes the required velocity is less than that the motor can reach
- Suitable gears can be used (adaptors) that diminish ω by a suitable constant factor. The ensemble motor + adaptor can be defined as *motoadaptor*
- The adaptors can be used individually or coupled to regulate the overall angular velocity

$$\omega_2 R_2 = \omega_1 R_1$$



Let's consider the mechanical power P conservation, $\Gamma \omega = \Gamma' \omega'$

Thus when the angular velocity drops down due to the effect of an adaptor the exerted torque correspondingly augments

In other words, if $\omega = n\omega'$ then $\Gamma' = n\Gamma$

Finally, since the mechanical behavior of a dc motor can be described by

 $\Gamma = \Gamma_0 + \gamma \omega + J \omega$ Γ_0 effective torque, γ friction and J inertia

 $\Gamma' - n\Gamma_0 = (J_R + n^2 J)\omega' + (\gamma_R + n^2 \gamma)\omega'$ Γ' new torque, γ_R friction and J_R inertia of the adaptor

Significant increase of inertia and friction Instabilities and hysteretical behavior due to mechanical play The motoadaptors can be effective where high precision is not required





78

Model of the mechanical part:

$$\tau = J_m \cdot \dot{\theta}_m + B_m \cdot \dot{\theta}_m + d + \tau_c$$

Similar approach to describe the load components:

$$\tau' = J_c \cdot \ddot{\theta} + B_c \cdot \dot{\theta} + d'$$

d = not manipulable inputs, frictions, phenomena without models

 $\tau_{c} = \text{effective torque applied to load} \\ \text{plus adaptor}$

 τ' = load torque







Model of the motor with load and adaptor





Adaptor hysteresis (due to possible mechanical plays)



DC motor, simplified equivalent electric circuit:

- • V_a = supply voltage
- • I_a = supply current
- •E = back electromotive force
- • R_a = armature resistance
- • L_a = armature inductance



Model:

- $\bullet V_a = E + R_a I_a$
- •E = $k_e \omega$
- $\bullet P_a = V_a I_a = E I_a + R_a I_a^2$

E back emf

 $k_{\rm e}$ constant for the velocity

absorbed electric power = mechanical power + power lost in the armature

•DC motor, characteristic:

- P_m = mechanical power = E $I_a = \Gamma_m \omega$
- Γ_{m} = motor torque = $k_{t}I_{a}$

•considering that:

```
I_a = (V_a - E)/R_a and E = k_e \omega we obtain that \Gamma_m = k_t (V_a - k_e \omega)/R_a
```



The torque diminishes when the angular velocity increases

At the beginning the mechanical power grows up with the angular velocity then diminishes with a typical "bell" behavior

TORQUE-VELOCITY REAL STATIC CHARACTERISTIC



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•DC motors, power supply:

Motors absorb a lot of current (0.5 : 1) A so requiring specific driving circuits (power drivers)

Two possible kinds of driving: linear and on-off



To manage a bi-directional rotation a bipolar power supply V_{in} +/- is required

If the resistant torque (the ensemble that opposes to rotation inertia, frictions, disturbs, ...) is negligible a good linearity can be obtained between $V_{in} \in \omega$

Should not a bipolar supply available, a "bridge" solution can be introduced

•DC motors, linear driving (a bipolar voltage using a unipolar one):



•DC motor, on off driving:

If the linear driving dissipates too much power, it can be used only in case of motors that absorb few Watts

To reduce the dissipated power MOS switching devices in on (saturation)/off stage



Dissipated power = V_{ce} I

If MOS off, I=0 no dissipated power

If MOS in saturation, $V_{ce}\,{\cong}\,0$ so dissipated power ${\cong}~0$

"Outflow" diodes for the current are required to avoid excessive voltages at the MOS drain (although it is fast and able to carry high currents)

Limit: the motor will run at the maximum velocity (only mono-direction) otherwise it will be stopped



•CC motor, on off driving:

For this reason it is useful to drive the motor through duty cycle variable pulses (PWM) so as to obtain a enough pulse frequency with an average supply voltage proportional to the pulses duty cycle.

However it is still not possible a bi-directional rotation



•DC motor, H bridge driving (for small laboratory robots):

the "H-bridge" circuit allows to control the sense of rotation.

Q1 e Q4 on	Clockwise rotation	V _{cc}
Q2 e Q3 on	Counterclockwise rot.	
Q1 e Q2 on	Q1 e Q2 can fire	
Q1 e Q3 on	Braking	
All the transistors off	Uncontrolled slowing down	

If two transistors on the same side are active (shoot-through) a big current passes between Vcc and ground. However, now the load is no more the motor (who absorbs the most of the current and so the voltage drop) as before thus the bridge could be damaged.

•CC motor, H bridge driving:

If we do not want to make the microprocessor serving as pulse emitter, a suitable electronic net could be used to activate transistors

 $V_{in} > V_{ds} Q_3 e Q_2$ active

 $V_{in} < V_{ds} Q_1 e Q_4$ active



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Output Spline

Servo Case

PWM PULSE EMISSION: SERVO MOTORS

They encapsulate dc motor, adaptors, control electronic e position control through a potentiometer for detecting the shaft position and rotation limit switch (to protect the movement area)

Characteristics: torque, rotation velocity, maximum rotation angle. Used in robotics applications

Motor

Drive Gears

Orders of magnitude:

Potentiomete

Control Circuit

Rotation angle: ±90°, torque: 3-20 Kg*cm (but even 20 Nm), absorbed current: 1-10 mA on rest and ~ 1 A full load, rotation velocity: up to 9000 rpm, weight: ~ 150 gr.

Input

GND

 $V_{cc} \cong 5 \div 6 V$



PWM PULSE EMISSION: SERVOMOTORS

•Servomotors:

PCM pulses (Pulse Code Modulation) are used, and the duty cycle determines the shaft position with respect the previous one (the incremental amount of the rotation): for example, depending on the motor type, a certain pulse duration corresponds to - 90° rotation, if the duration is twice the rotation is +90°, if it is half the rotation is 0°. If the servo does not frequently receive pulses it is abandoned without any control.

Usually the repetition period is 20-30 msec, whilst the pulse duration ranges from a 1 msec minimum up to 2 msec max.

The servo is stopped when a "central" pulse duration is provided or due to its natural lost of energy (frictions ...)

Open loop position control



PWM PULSE EMISSION: SERVOMOTORS

•Servomotors:

The "response" of the servo depends mainly on the pulse duration more than on the duty cycle since these motors are designed to properly work into a enough wide pulse frequency range within which the behavior is quite similar (not a orthodox PWM).

If for example a servo receives pulses with variable period among 0-70 msec, it may happen that with values close to ~10 msec the servo becomes a little unstable (vibrations), while when the period is > 40 msec, the torque diminishes. Thus a period close to 20 msec could be the right choice (moreover is easy \Rightarrow 50 Hz).

The servo are characterized by a suitable "working curve" to achieve a predetermined motion dynamic.

The "curve" could consist on variable width steps (larger at the beginning, smaller when close to the target) up to reach the set point. It could be useful to control the movement velocity although this is quite difficult since practically the system moves always at maximum velocity.

How to generate the pulse?: a hybrid solution could envisage a μ P triggering a timer so as it exits pulses with a duty cycle set again through the μ P. It is not possible to completely devote a μ P to pulse emission because the required times (sometimes < msec) ask for high frequency routines, with possible μ P overload and inaccuracies.



- •At the beginning the step is wide, when the position is close to the target it diminishes
- •The position of the rotating shaft can be evaluated time by time

- •Step motors: every received pulse correspond to a precise rotation step
- •Low torque, good precision, low cost. Applications: computing systems peripherals, robotics, small applied loads

• Specs:

- •Pulse duration enough to drive electronics
- •Start frequency (minimum sequence necessary to start rotation) and stop (minimum frequency below which the motor is stopped)
- Acceleration and deceleration ramp
- Possible jitters that imply steps lost or vibrations



•Dedicated hw or processors (for example PIC: programmable interface controller, Microchip Risc microcontroller, low costs)





Most commonly used movement (good torque, not negligible power consumption and heat dissipation)

1	2	3	4	V ₁₃	V ₂₄
ON	ON	OFF	OFF	+	+
OFF	ON	ON	OFF	-	+
OFF	OFF	ON	ON	-	-
ON	OFF	OFF	ON	+	-
ON	ON	OFF	OFF	+	+



How to provide clock?

1	2	3	4
ON	OFF	OFF	OFF
OFF	ON	OFF	OFF
OFF	OFF	ON	OFF
OFF	OFF	OFF	ON
ON	OFF	OFF	OFF

Wave: less torque with respect two phase (1vs1.4)

<i>Half step</i> : the step	1	2
number is twice	ON	OF
(precision) but the	ON	OF
torque is irregular	ON	O
and the power	OFF	O
consumption not	OEE	
constant		U

1	2	3	4
ON	OFF	OFF	ON
ON	OFF	OFF	OFF
ON	ON	OFF	OFF
OFF	ON	OFF	OFF
OFF	ON	ON	OFF
OFF	OFF	ON	OFF





Fig3 - One-two phase on - half step





1a. WAVE DRIVE M	ODE	1b. NORMAL DRIVI	E MODE	1c. HALF STEP MODE		LEB WODE	
Transistor Accesi	Posizione Rotore	Transistor Accesi	Posizione Rotore	Transistor Accesi	Posizione Rotore	Transistor Accesi	
Q1	s N	Q1 e Q2	S N N S	Q1	N S S	Q3	
Q2	s(n s)n→	Q2 e Q3	S N N N	Q1 e Q2	S N N S	Q3 e Q4	
Q3	s N N	Q3 e Q4	N S N S	Q2	s(N S)N→	Q4	
Q4	N(S N) S 🔶	Q4 e Q1	N S S	Q2 e Q3	s N N N	Q4 e Q1	

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Posizione Rotore s

Ν

Ν



HSTEP motorSet motor to use half-step drive sequence.
2PHASE motorSet motor to use two-phase drive sequence. This is the default drive sequence.
WAVE motorSet motor to wave drive sequence
RATE motor val1 <val2>Set the delay rate between steps (in milliseconds). Range from 1 to $50,000$.</val2>
DIR motor <cw, ccw="">Set the direction of rotation. CW = clockwise, CCW = counter-clockwise. If no direction is given then it is reversed.</cw,>
STEP motor val1 <val2>Step motor by val1 steps. If motor is both then val1 refers to `X" and val2 refers to `Y'.</val2>
WAIT motorWait for the previous motor command to finish before executing the next command for that motor.
SPIN motorContinuously step motor. Not affected by `WAIT'.
STOP <motor>Immediately stop the current command for motor. If no motor is specified then stops both. Ignores 'WAIT'.</motor>
MANUALManually step motors. Another menu appears.
LOOPTIL input(n) HIGH,LOW where n=1, 2, 3 or 4.Wait for the specified input to go high or low before continuing.
DELAY millisecondsWaits the specified delay time before continuing. Range = 1 to 50,000.
ECHO ON or OFFIf ON (default) then all commands are echoed to the output.
PRINT messageto output
VERPrint version number
ENDWait for all motor commands to finish then quit
QUITQuit program immediately
HELPDisplay above summary

STEPPER MOTORS: how to power supply

•This actuator responds with a velocity to the control input variable, allowing to determine the new position without any need of a position transducer (i. e. if we fix a 100 pulses/sec velocity, this corresponds to half a round in a second, if the sensibility of the transducer is 1 round after 200 pulses). Sometimes a limit switch is necessary.

•These motors require power supply in the order of few V, with currents in the order of the A \Rightarrow r in the coils must be relatively low; since the delay is L/r in an inductor, time constants can be high. L moreover can be high if the coils windings are high.

For example if the command voltage is applied to the coil 1, this one due to the initial 'resistance' will inhibit the, at least at the beginning, the current flow. This means that in the MOS drain V_{cc} is nearly present and that the MOS will employ further time to exit from the interdiction towards saturation \Rightarrow this is acceptable if the working frequency is not so high.

It is better to provide higher voltages (\Rightarrow > resistors, < time constants): however this brings to a strong power dispersion on the MOS (heat). Thus a *PWM* solution is employed with suitable on-off intervals (tunable duty cycle).

STEPPER MOTORS: PWM POWER SUPPLY



The motor receives a constant current, obtained through a HV voltage necessary to provide fast commutations (time constant L/r)

If current is $< V_{rif}/R$ the flip flop close the switch so supplying the motor with a 100 V voltage. Alternatively (open switch), the circuit is grounded through a diode.

The flip flop closes the switch depending on a clock period that allows to regulate the opening and closing intervals (that means the duty cycle).

The inductor serves as "regulation maker" allowing a smooth passage among the levels of the current (high low high low) without shocking the motor.

ON OFF control.



- "controlled" AC power supply through SCR/TRIAC
- The wave passage through zero must be revealed then a μP should provide the pulses for SCR/TRIAC activation
- Wave period 20 msec, the passage through zero has to be detected with very high precision \Rightarrow HW solution
- D(t) can be settled through a sw routine or through a timer with a interrupt to μP
- Non linear relationships between the average power provided and the power up delay



 $\frac{2}{3}Vcc$



$$Ic = \frac{Vcc - Vb - 0.7}{R}$$

 $\mathbf{T} = \frac{2 \ RC \ Vcc}{3 \ (Vcc - Vb - 0.7)}$

 V_b allows to regulate the capacitor charge velocity so to establish when (τ) to power up the SCR (i. e. the provided power).

High non linearity between the delay and the power applied to the load.



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CODED SET POINT VALUES



HW CODED (SETUP) VALUES: BCD SWITCHES

Signals alone do not bring significant information: it's their ordered ensemble that represents an information

A typical example is the acquisition of BCD numeric values for the digital set point of a regulation chain = *contraves*



HW CODED (SETUP) VALUES: BCD SWITCHES



OEM COURTESY

Possible problems during commutations:

- •Brushing, slithering and contacts rebounds
- •Irregular commutations due to long usage
- Passage through halfway not significant values
- •Casual order through which the operator moves digits
- •Several input instructions (if for example more than 4 BCD digits and μ P with <= 16 bits)

Solutions:

- •Ok temporarily "out of range" values
- •A suitable "settlement time" (SW) after which values are accepted
- •Validation button (HW)

•A sw conversion from the BCD value to the machine binary representation is needed ...



CODED VALUES: ABSOLUTE ENCODERS

Gray code is used to minimise the error due to a bad alignment of the optics internal to the encoder.

Pseudocode algorithm

```
commutare := 0
i := N-1
while i >= 0
BitBin[i] := BitGray[i]
if commutare = 1
then BitBin[i] := not BitBin[i]
commutare := BitBin[i]
i := i-1
endwhile
```

C version

```
int Gray12_bin (NumGray);
{
    int commuta,i,mask;
    valore = NumGray;
    i = 12;
    mask = 0x800;
    commuta = 0;
    while (i > 0)
    {
        valore ^= commuta;
        commuta = valore & mask;
        commuta >>= 1;
        mask >>= 1;
        i--;
    }
    return (valore);
```

ASSEMBLY VERSION

Routine that takes as input the variable VALORE in Gray code 12 bit and exits with the pure binary converted value in AX (accumulator) register.

,	Integer	Gray code	Binary code
, MOV CX,12 MOV AX,80 SUB BX,82 CICLO: XOR [VALC MOV BX,[V AND BX,A2 SHR AX,1 SHR BX,1 LOOP CICLO MOV AX,[V RET	Integer 0 0 1 0 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 <td>0000 0001 0011 0010 0110 0110 0111 0101 0100 1100 1100 1101 1111 1111 1110 1010 1011 1001 1000</td> <td>Binary code 0000 0001 0010 0011 0100 0101 0101 0111 1000 1011 1001 1011 1100 1101 1101 1111</td>	0000 0001 0011 0010 0110 0110 0111 0101 0100 1100 1100 1101 1111 1111 1110 1010 1011 1001 1000	Binary code 0000 0001 0010 0011 0100 0101 0101 0111 1000 1011 1001 1011 1100 1101 1101 1111

OUTPUT CODED VALUES



We suppose that we have to output more bits than those manageable with a single OUT instruction

OUT (PIOA), reg OUT (PIOB), reg1 OUT (PIOC), WR_CMD

DATA BUS ADDRESS BUS