# INDUSTRIAL INFORMATICS 

## and EMBEDDED SYSTEMS

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## COURSE SECTIONS:

- DIGITAL INTERFACING
- INDUSTRIAL COMMUNICATIONS \& FIELD BUSES
- EMBEDDED ARM MICROPROCESSORS


## THE EXAM

Typically: 1) a code fragment in ARM assembly language to be interpreted<br>2) an exercise on the evaluation of the frequency of the pulses provided by a suitable sensor, considering a microprocessor whose bits are employed for different purposes (driving a motor, acquisition of a physical a magnitude, generating a signal, PWM, PFM, ...)

As the case is: 1) questions on the different parts of the course to establish if mechanisms have been really captured
2) exercises on the communication protocol part

## MASTER THESIS (finalised to job positions)

Companies that propose thesis:<br>MARELLI MOTORSPORTS (Corbetta Milano): video data logging, telemetry<br>TEMIS (Corbetta Milano): automotive, satellites, data acquisition and communication<br>BDSOUND (Assago): sound engineering, audio signals acquisition through microcontrollers (Cortex M4, ST, ...)<br>AZCOM (Rozzano): wireless (and not) safe communications through sw control (DSP and FPGA processors)<br>FACILITY LIFE: cybersercurity<br>ST MICROELECTRONICS: embedded processors

## MASTER THESIS (in laboratory)

## DSP, GPU, FPGA technologies applied to:

Brain cerebellar simulation
Hyperspectral image processing for high accuracy cancer detection Glaucoma identification in ophtalmic imaging

Automotive applications
Signal processing through embedded microprocessors

## FOCUS OF THE COURSE


-Control algorithm = the controller operating principle
-Process = the problem. Know the transfer function - Minimize the gap between SP and PV (and dependency on environmental variables)
-Oscillations due to the velocity of the answer to SP variations $\Rightarrow$ minimize the transient
-Overshooting $\Rightarrow$ limit the amplitude
-Stability (no infinite or diverging oscillations)

## COURSE SECTIONS

- DIGITAL INTERFACING
- INDUSTRIAL COMMUNICATIONS \& FIELD BUSES
- EMBEDDED ARM MICROPROCESSORS


## DIGITAL INTERFACING

Interfacing = interactions between active operative devices (i. e. CPU) and/or passive ones (i. e. memories, sensors, ports)

The course is mainly (but not only) focused on the interactions between calculator and external world (every thing not directly manageable by the CPU through buses)

Several kinds of in-homogeneities:

- Physical: the external world is characterized by physical parameters that can be represented only through electric entities with suitable amplitudes (transducers)
- Formal: codification (BCD/Gray), representation (pulse count or frequency measurement), signal processing/management (PWM or PFM modulation)
- Temporal: how external phenomena are synchronized with signal acquisition/emission? (timers, interrupt/polling/DMA, real time protocols)
- Spatial: computers and controlled devices are not physically adjacent (transmission)
- Attitudinal: Human Computer Interaction

A suitable chain of interacting devices allows to overtake these in-homogeneities

## DIGITAL INTERFACING

## Typical computer-external world interfacing

Interface design = identification of a chain of components able to wholly satisfy the requirements of data transformation and communication


BUS


## A general interface digital circuit

-I/O buffer (Fifo)

- State registers for significant situations (errors, overrun, interrupt ...)
- Command registers to setup mode functions and modify state bits
- Various functions. DMA, ser/par transfers


## ADDRESSING ON THE BASIS OF A MEMORY MAP



## ADDRESSING ON THE BASIS OF A MEMORY MAP

1) A single address space (MEMORY MAPPED I/O)
lw \$reg1, off(\$reg)
\$reg= port (mem) address
ex. ARM, MIPS
2) Double address space (CUSTOM INSTRUCTIONS)

Iw \$reg1, off(\$reg)
in \$dest, (\$reg)
ex. INTEL
\$reg=memory address
\$reg=port address

Control signals to distinguish between addressable space devoted to data

memory and I/O access (MREQ, IORQ) addressable (duplicated) space devoted to peripherals


## DIGITAL INTERFACING

## On off signals $\quad$,

-They represent physical entities acquired through single bits of a microprocessor port -What is important is the value of the bit, not its variation
-8-16-32 bit groups acquisition (bytes, words)
-Let's test the Kth bit

```
Esempio in C
    var_bool = 0;
    if (immagine & mask_K)
        var_bool = 1;
Variante più compatta:
    var_bool = (immagine & mask_K) && TRUE;
```

-And now two assembly routines (the right one is optimized maybe by the compiler)


## DIGITAL INTERFACING

## On off signals

-Let's set the Kth bit

```
Esempio in C
if (var_bool)
    immaḡine |= mask_K; // OR con mask
else
    immagine &= ~ mask_K; // AND con COMPLEMENTO di mask
```

```
Esempio in Assembler
    MOV AL,MASK_K ; maschera
    AND AL, [VAR_BOOL] ; isola bit in posizione K
    MOV AH, AL ; salva in AH
    MOV AL,MASK_K
    CPL AL ; maschera complementata
    AND AL, [IMMAGINE] ; azzera bit in posizione K
    OR AL, AH ; inserisce nuovo bit in posizione K
    MOV [IMMAGINE], AL ; aggiorna variabile immagine
```


## DIGITAL INTERFACING

## On off signals: "state" acquisition

- The value of a bit is read to establish the level (state) of a line related to a physical entity: use of a three state buffer



## DIGITAL INTERFACING

## On off signals

- three state buffers to de-couple the line with respect other devices



## DIGITAL INTERFACING

## COTS buffer components (Texas Instr. SN 74AC244)



SN54AC244 . . FK PACKAGE (TOP VIEW)

logic diagram (positive logic)


## DIGITAL INTERFACING

## On off signals: event triggering

-The value of a bit is read just to establish if a commutation has occurred (event) of a line related to a physical situation, by comparing it with the previous value
-2 possible approaches:
-SW triggering through polling
>A microprocessor port is continuously sampled. The software compares the new value with the old one and identifies the eventual commutation
-triggering through interrupt
>The signal carrying the "event" is connected to a microprocessor port's pin
$>$ The single commutation edge is determined through pins edge triggering (falling/raising edges)
>Both the edges are acquired through suitable electronic circuits sensible to edge or to levels. Where is the difference?

## DIGITAL INTERFACING

## On off signals: "event" acquisition through interrupt

Latch (sensible to the level)
FF (sensible to the edge)


## DIGITAL INTERFACING

## On off signals

- a bit value is read to establish the possible commutation using one-edge sensible circuits. Below a circuit able to transform commutation of both the edges in a single direction

$T_{c k}<T_{s} / 2$ if duty cycle 50\%


## DIGITAL INTERFACING

## On off signals

a bit value is read to establish the possible commutation using level sensible circuits (a FFJK, or a SR, is required to keep the input fixed)


## DIGITAL INTERFACING

## Physical signal adjustment: opto-couplers

- Mandatory if grounds are different or if possible over-voltages
$\bullet R_{\mathrm{L}}$ to determine necessary current for LED power up
$\bullet R_{p}$ pull-up (saturation vs "steep" edges: what possible values?)



## DIGITAL INTERFACING: FILTERING

## On off signals

-This kind of signals are produced by relays, electronic switches, buttons, limit-switches
-Possible spurious states irrelevant if the signal level is required but misleading if we are interested in the commutation (i. e. an associated event)

-electric vs algorithmic filtering
-electric. required if interrupt o counters are used ( $\Rightarrow R C+$ Schmitt trigger)
-algorithmic. necessary if the signal is sampled. The sampling will turn out in a series of values among which some are significant others not (spurious). The choice among them is driven by the information that the signal carries $\Rightarrow$ definition of a suitable transient death time

## DIGITAL INTERFACING

Fine corsa a leva, stelo, pulsante
Lever, stem, push-button,

Fine corsa ad albero
tree limit switches


## DIGITAL INTERFACING

## Electric filter through RC circuit + Schmitt trigger

Choose carefully the values of $R$ and $C$, since they influence the circuit time constant and thus the reliability of the microprocessor acquisition

Slow capacitor charge/discharge could imply a slow crossing in the uncertainty range, bringing an oscillating behaviour. Therefore a Schmitt trigger could be suitable.


Simple filtering circuit for input signal to be read by a $\mu \mathrm{P}$


Input waveforms before and after filtering

## DIGITAL INTERFACING

## Schmitt Trigger (Filter)

-Its characteristic is an hysteresis
-It depends on the behavior of a magnitude (memory effect): different output according to the direction of variation
-Useful in case of slow spurious signals and hw detection of a commutation
-The input signal is "squared" and is more suitably distinguished by the port electronic circuits


## DIGITAL INTERFACING

## Algorithmic filtering: first edge detection

-The first variation is considered as really significant and accepted
-Further variations are ignored during the time interval T (transition death or dead time)
-The algorithm does not work with isolated spurious signals
-Spurious signals follow a real significant event and run out within the T interval. Ok in case of mechanical contacts


## DIGITAL INTERFACING

## First edge detection, a simple assembly code

```
; Routine da chiamare ciclicamente ad ogni periodo di campionamento.
; Riporta in AL il valore filtrato
ACQ_BIT_N:
    CMP [TRANSITORIO],0
    JE REGIME
;qui in transitorio - conta il tempo trascorso
    DEC [TRANSITORIO]
    JMP ESCI
;qui transitorio esaurito - pronto a sentire event. commutaz.
REGIME:
```

IN AL.PORTA IN
AND AL,MASK_N
CMP AL,[PRECEDENTE]
JE ESCI
;qui rilevata commutazione
MOV [PRECEDENTE],AL ;aggiorna il valore
MOV [TRANSITORIO],T_MORTO ;rilancia il transitorio
ESCI:
MOV AL,[PRECEDENTE] ;riporta il valore valido

## DIGITAL INTERFACING

## Algorithmic filtering: expired transient detection

-The first variation is not significant
-The corresponding state is accepted as valid only if unchanged within T time interval
-The algorithm works well with isolated spurious signals
-A low pass filtering effect (delay)


## DIGITAL INTERFACING

## Exhausted transient detection, a simple assembly code



## DIGITAL INTERFACING

## On off signals emission

-The output on a $\mu \mathrm{P}$ port is carried out through latches (permanent values)
-Output of initial setup values or after a reset (0 if possible)
-Or simply those values that will manipulate the process through the actuator
-Output through "image" variables (it is not possible to output single bits)
-Image variables processed through AND/OR/XOR masks
-the National Instruments Latch 54ACT374 component

## DIGITAL INTERFACING

\section*{Pin Assignment for DIP and Flatpak <br> 

Logic Diagram


Please note that this diagram is provided conly for the understanding of logic operations and should not be used to estimate propagation delays.

Pin Assignment for LCC


Truth Table

| Inputs |  |  | Outputs |
| :---: | :---: | :---: | :---: |
| $D_{n}$ | $C P$ | $\overline{O E}$ | $O_{n}$ |
| $H$ | $\sim$ | $L$ | $H$ |
| $L$ | $\sim$ | $L$ | $L$ |
| $X$ | $X$ | $H$ | $Z$ |

$\mathrm{H}=\mathrm{HIGH}$ Vollage Level
L = LOW Vdlage Level
$\mathrm{X}=$ Inrmaterial
Z = High Impedance
$\sim=$ LOW-b-HIGH Transition

Latch 54ACT374
National Instruments

## DIGITAL INTERFACING

## Pulses


Binary digital signal
commutations to which a
single event is associated
considered as significant

Sometimes both the edges are useful since associated to the beginning and the end of a event $\Rightarrow$ interval measurement

## Pulses can be acquired for:

-Temporization
-Counting a suitable number of moving objects
-Calculating object position
-Calculating object velocity

## Pulses can be emitted for:

-PFM device driving
-PWM device driving
-Power supply to an electric load through partialization of the signal

## DIGITAL INTERFACING

## Pulse acquisition for counting (object passage, complete rounds ...)


-LED + phototransistor to reveal the object passage
-Total primary counter must be upgraded (i. e. program variable)
-Max reachable value => type definition => unsigned long ( 32 bit)
-Non volatile counting (eventually secondary counters in powered buffer memories / files. Trade-off between freshness and CPU overload)
-Manual vs automatic reset operations
-Variables must be stored through atomic saving operations in memory

## DIGITAL INTERFACING

## Hw primary counter

-Less sw elaboration charge onto the CPU
-Mandatory if pulse frequency not acquirable through polling from the $\mu \mathrm{P}$ port (it requires few assembly instructions)

- If the counter is connected with a $\mu \mathrm{P}$ port we read a state information
-If the counter is connected with an interrupt line (triggers the interrupt) we read an event information



## DIGITAL INTERFACING

Hw primary counter: a commercial component Texas SN74193

-Clear pin to reset
-Up/Down pins for setting bidirectional inputs
-Load pin to values preset

- $A_{r}$ B, C, D preset inputs
-Borrow/carry pins



## DIGITAL INTERFACING

## External (primary) hw counter providing a "status" information

N bit hw counter

- $\mu$ processor port connection
-Acquisition routine periodic activation $\mathbf{T}_{\mathbf{c}}$
where $\mathrm{T}_{\mathrm{c}} \leq\left(\mathbf{2}^{\mathrm{N}} \mathbf{- 1}\right) * \mathrm{~T}_{\text {imp }}$

Input from the port (present counter value)
\# pulses between two inputs ( $\mathrm{V}_{\text {new }}-\mathrm{V}_{\text {old }}$ added to a total (secondary) counter (carry management)

AGGIORNA CONTATORE:
SUB AX,AX
IN AL, PORTA CONT SUB AL, [PRECEDENTE] ADD [PRECEDENTE],AL ADD [CONTATORE],AX ADC [CONTATORE + 2],0 RET

## DIGITAL INTERFACING

## External (primary) hw counter providing a "status" information

If $\mathbf{n}^{\circ}$ of bits > than readable by a single instruction
-Bytes from the counter are read with successive read phases
-Auxiliary register with the same width of the counter to be read in two phases


AGGIORNA CONTATORE:
IN AL,P_MSBYTE
RIPETI:
MOV BH,AL
IN AL,P_LSBYTE
MOV BL,AL
IN AL,P_MSBYTE
CMP AL,BH
JNE RIPETI
;qui $\mathrm{BX}=$ valore valido
MOV AX,BX
SUB AX,[PRECEDENTE]
ADD [PRECEDENTE],AX
ADD [CONTATORE],AX
ADC [CONTATORE + 2],0

## DIGITAL INTERFACING

## External (primary) hw counter providing an "event" information

-Interrupt line connected to a single counter bit or to the carry bit -Used as "pre-scaler" (frequency division)

$$
\begin{gathered}
F_{\text {out }}=F_{\text {in }} / 2^{\mathrm{N}}(\text { if carry bit connected }) \\
F_{\text {out }}=F_{\text {in }} / 2^{\mathrm{k}}\left(\text { if } k^{\text {th }} \text { bit connected }\right)
\end{gathered}
$$

-Memory SW counter upgrading (if high frequencies)
-Least significant count directly read from the counter

## SW counter

 Direct acquisition
## DIGITAL INTERFACING

## Pulses to evaluate an object position

-Pulses corresponds to incremental movements. Where is the "zero" point?
-Is it an un-surmountable extreme? Unsigned numbers
-Is it an arbitrary position? Signed numbers
-The incremental encoder

A wheel made up by windows alternatively transparent and opaque to a light beam emitted by a LED source

The light is received by a phototransistor

Each window provides a binary information (i. e. a bit)


## DIGITAL INTERFACING

## Incremental Encoder

-Two concentric crowns, $\mathbf{1 / 4}$ period out of phase
-Two couples of LED/transistor (1 light, 0 dark).
-Minimum measurable angle = half window


The two acquired signals (pulse trains) as they


## DIGITAL INTERFACING

## Incremental encoder



## DIGITAL INTERFACING

## Counting specs

- definition of "zero" (natural/integer counting variable)
- maximum counting vs granularity (size)


## Counting

- edge detection
- movement direction
- primary counter upgrading (then eventually upgrade the secondary one)
- 2 possible approaches: sw/hw


| A | 13 | C |
| :---: | :---: | :---: |
| P | O | - 1 |
| P | 1 | -1 |
| N | O | -1 |
| N | 1 | $\rightarrow 1$ |
| 0 | P | -1 |
| 1 | P | $\rightarrow 1$ |
| 0 | N | $\rightarrow 1$ |
| 1 | N | -1 |

## VELOCITY CALCULATION THROUGH AN ENCODER - sw approach

-A and B: 2 pulse trains connected to two pins of the $\mu$ P port
-Commutations detection
-Movement direction identification: forward (+) or backward (-)
-Counter variables (primary and eventually secondary)
-algorithm (hypothesis clockwise rotation considered as positive)
if $B$ switched
if $B=A$ count ++
else count--

else count + +
end if

## DIGITAL INTERFACING

## An example of sw management of a bi-directional incremental encoder



## Routine activation frequency?



We have to distinguish between:
case ok $B(t-1) \neq B(t)$ with $A=1$, (cw)
case ko $B(t-1) \neq B(t)$ with $A=0(c c w)$

$$
F_{c l k}=2 * F_{s}
$$

This on both the A e B signals
Fclk $=4 *$ Fs

## DIGITAL INTERFACING

## Hw management of a bi-directional incremental encoder



The hw logic must distinguish among commutations

Activation frequency? of the same signal (A o B); however, within the present commutation and the next one it must also recognise if the other signal is $\mathbf{0}$ or 1.

$$
F_{\mathrm{clk}}=8 * F_{\mathrm{s}}
$$

## DIGITAL INTERFACING

Hw management of a bi-directional incremental encoder


## DIGITAL INTERFACING

## Hw management of a bi-directional incremental encoder



Only with with $8 \mathrm{~F}_{\mathrm{s}}$ sampling frequency it is possible to reveal useful commutations and the turnover among useful configurations and "sleep" states

Hw management of a bi-directional incremental encoder

CW rotation


Whot hoppena iff $F_{\text {cik }}$ is $4 F_{\text {ssumare }}$

Hw management of a bi-directional incremental encoder

CCW rotation


CW rotation
whot hoppena if $F_{\text {cik }}$ is $4 F_{\text {suminc }}$

$$
\begin{aligned}
& \text { DEC } 2 \quad 4 f_{5} \int_{2^{\circ} \text { asso }}^{3^{00450} 0,2,7,7,4} 0
\end{aligned}
$$

## DIGITAL INTERFACING

## Pulses for timing

Time-stamping means to associate events and times into which they happen
Errors due to the model describing time evolution within the machine (TIG - temporal granularity)
Delays due to timer acquisition sw ( $\mathrm{T}_{\mathrm{e}}$ ) and to event recognition ( $\mathrm{T}_{\text {rev }}$ )
Latencies = time interval between the event and its perception ( $\mathrm{T}_{\mathrm{lat}}$ ). Non deterministic, if they are due to resources charged with task perception and in that moment devoted to other activities (priority). Typical in interrupt service and in processes activation


## DIGITAL INTERFACING

## Pulses for timing (maximum priority routine)

Events internal to the machine (i. e. switch of a control unit signal)

$$
\mathrm{t}-\mathrm{TIG}+\mathrm{T}_{\mathrm{el}}<\mathrm{T}_{\mathrm{st}}<\mathrm{t}+\mathrm{T}_{\mathrm{el}}
$$

Events external to the machine (i. e. switch of a signal read through a port)

$$
\mathrm{t}-\mathrm{TIG}+\mathrm{T}_{\text {rev }}+\mathrm{T}_{\mathrm{el}}<\mathrm{T}_{\mathrm{st}}<\mathrm{t}+\mathrm{T}_{\text {lat.max }}+\mathrm{T}_{\mathrm{el}}+\mathrm{T}_{\text {rev }}
$$

Measurement of an interval between external events (constants components neutralise each other, deterministic ones do not disappear)

$$
- \text { TIG }-T_{\text {lat.max }}<\left(\mathrm{T}_{\mathrm{st1} 1}-\mathrm{T}_{\mathrm{st2} 2}\right)<\mathrm{TIG}+\mathrm{T}_{\text {lat.max }}
$$

What is TIG? It depends on the machine time representation model. Let us open a parenthesis ...

## TIME CHARACTERISATION

- Is the time an event container (box) or the events cause the course of the time?
- During life we follow the first approach, in computer science we are close to the second one
- CLOCKS = phenomena that generate regular events (ticks) used to quantify the time passing
- We need a model


## TIME AS IT IS REPRESENTED WITHIN A CPU

-Hp: bijective correspondence among time moments and real numbers
-Continuity, limit, derivative, integral existence ecc.
-Time as linear, monotonic, discrete
-Clock $=$ time events generator with domain $\mathbf{N}^{+}$
-Quantization = duration of the phenomenon taken as clock, usually very negligible with respect events that must be measured

## TIMING MODEL

-LOCAL TIME UNIT (UTL) = absolute time interval corresponding to 1 in the measurement unit of the clock
-LOCAL TIME BASE (BTL) = multiple of the UTL, representing the period of the cyclic phenomenon assumed as clock
-Let' suppose that in a PC the real time clock is 55 msec: the time is represented in msec (UTL), but it is increased every 55 (BTL)
-UTL = resolution, BTL = granularity
-If $\operatorname{BTL}=1$, resolution = granularity
-If T = clock time and t = absolute time,

$$
\begin{gathered}
\text { T}=\mathbf{n} * \mathrm{BTL} \\
\mathrm{n} * \mathrm{BTL} * \mathrm{UTL} \leq \mathrm{t} \leq(\mathrm{n}+\mathbf{1}) * \mathrm{BTL} * \mathrm{UTL}
\end{gathered}
$$



## TIME MODEL: PRECISION ERRORS

-They are due to how the phenomenon (assumed as clock) differs from the ideal behavior
-PRECISION
A clock is precise if $t(t i c k(n))-t(t i c k(n-1))=B T L * u t l+e(n)$ with $e(n)=0$

- Systematic error $\Rightarrow$ the phenomenon taken as clock features a not correct frequency
- Null mean errors $\Rightarrow$ not stable frequency (jitter)
- precision measurement $=$ BTL*utl $/$ mean_err $=\left(\mathbf{1 0}^{4}-10^{6}\right)$
-CORRECTNESS
A clock is correct if abs(T-t/UTL) < BTL, that is the local time $T$ is synchronised with the absolute time (within the quantization error BTL)
-A precise clock if is correct in $t$, is correct in $t^{\prime}>t$
-Sistematic unaccuracy (mean_err $=0$ ) $\Rightarrow$ no correctness
-If intervals are measured, systematic errors do not exert any influence


## TIME MODEL: QUANTIZATION ERRORS

-They are caused by the representation of the time
-ORDER: it is not possible to establish an order between 2 events that happen within $\mathbf{2}$ consecutive ticks (i. e. within the same BTL)
-CONTEMPORANEITY: if the time interval between two events is $\Delta T<B T L$, the computer considers them as simultaneous (contemporaneous)

## -MEASUREMENT:

- 2 events, whose real time distance is $\Delta T$, are considered as within $\mathbf{N} * \mathbf{B T L}<\mathbf{M}(\Delta T)<(N+1) * B T L \quad N=$ floor $(\Delta T / B T L)$
- events, whose time distance is $N * B T L$, can correspond to a real time interval $\Delta \mathrm{T}: \quad(\mathbf{N}-\mathbf{1}) * B T L * U T L<\Delta T<(N+1) * B T L * U T L$



## TYPICAL TIMING SYSTEMS (TIMERS)

The aim is to have different granularities (1:100) one for time measurement (accurate) and/or one for process activation (coarse)


Oscillator ( $\mathrm{F}=20 \mathrm{MHz}$ )
Prescaler (frequency division, $2^{\mathrm{k}}=1-256$ )
TIMER $=\mathrm{N}$ bit binary cyclic counter (16-24)
TIME BASE COUNTER = 1 tick when the set period (TBV) ends + restoration TBV
Decrement counters

## TYPICAL TIMING SYSTEMS (TIMERS)

- TIMER section

TIMER RESOLUTION (count units) = 2k/F=UTL
TIMER PERIOD $=2^{N} * 2^{\mathrm{k}} / \mathrm{F}$
$B T L=2^{N}$ (on the overflow)

$(C O U N T+T B V) \bmod 2^{M}=C O U N T+T B V-2^{M}$
$=2^{M}-x+T B V-2^{M}=T B V-x$
to include counts lost due
to interrupt latency

## - TIME BASE section

TIME BASE COUNTER RESOLUTION = 2k/F=UTL

TIME BASE COUNTER PERIOD = TBV

* $2^{\mathrm{k}} / \mathrm{F}$

BTL=TBV

## TBV RESTORE

-Set TBV automatically or ...
-Explicit setup within the sw routine that manages tick interrupt but ...
-(CONTEGGIO+TBV) mod2²

## SOME IDEAS

```
F=1 MHz, Prescaler=1 = period 1 }\mu\textrm{sec},\mathrm{ if N=16 = 65000 }\mu\textrm{sec
F=1 MHz, Prescaler=256 => period 256 \musec, if N=16 => 16 sec
F=20 MHz, Prescaler=256 => period 12.5\musec, if N=24 => 200 sec
```

Process activation
Time stamping (date and time) $\quad \Rightarrow$ tick TBC
$\Rightarrow$ overflow TIMER
$\Rightarrow$ TIMER

NB. If long intervals are measured we need a sw variable onto which to store the counter values. This implies a partial un alignment between the total sw counter value and the hw register (containing something like a <fractional part» of the overall value).

## DIGITAL INTERFACING

## Pulses are acquired to calculate the velocity (i. e. frequency) of an object

Angular velocity or (scale factor) linear velocity can be determined
Problem data

- Encoder granularity = (round or meter/pulse) $)^{-1}=>$ quantization effect
- Maximum and minimum frequency to be measured
- Pulse (usually given by a sensor) duration


## Specs

- Measurement production frequency (acceptable depending on the particular phenomenon to be measured)
- Delay between the acquisition instant and the moment to which it is related
- Errors (relative vs. absolute, related to measurement value or to the full scale value, Erm\% vs. Erf\%)


## VELOCITY CALCULATION

Pulses are counted during a known time interval


The period between two pulses is evaluated Tck


What is measured is an "average velocity"

## Pulse acquisition techniques

## -Software

Cyclic (Tc period) pulse acquisition
If time for reading negligible and Tg granularity of the used clock (BTL* UTL) Error on the event detection: $-\mathrm{Tg}<\mathrm{E}<\mathrm{Tc}$
-Interrupt
The signal carrying the pulse is connected to the interrupt pin of a microprocessor
$\mathrm{T}_{\text {lat.max }}=$ max. latency of the interrupt service routine
$\mathrm{Te}=$ execution time of the interrupt service routine
Error on the event detection: $\quad-\mathrm{Tg}+\mathrm{Te}<\mathrm{E}<\mathrm{T}_{\text {lat.max }}+\mathrm{Te}$

## -Hardware

Dedicated circuits with working times negligible with respect to event duration

$$
\text { Error on the event detection: } \quad-\mathrm{Tg}<\mathrm{E}<0
$$

## Pulse count during time interval: I-SW, T-SW

-Sw routine executed with period Tc:

- Input port read
- if positive (raising up) edge I count++;
- if positive edge T vel=(enc. granularity)*counter_val/ $\Delta \mathrm{T}$; count=0;


## -Constraints

$\mathrm{Tc}<$ minimum pulse duration $\quad \Rightarrow \quad \mathrm{f} T \mathrm{c}<1 / 2$ ( $\mathrm{f}=$ pulse frequency)
$f_{p}$ (measure production freq.) $=1 / \Delta T ; \quad f_{\text {min }}$ (minimum detectable frequency) $>1 / \Delta T$

## -Errors

DTM $=$ real (effective) count interval; $n_{p}=$ real $n^{\circ}$ of pulses DTM
$f=$ real signal frequency $=n_{p} / D T M ; f_{m}=$ measured frequency $=n$ (counter value) $/ \Delta T$

Quantization err. on pulse identification [-1 ... 1] $\Rightarrow n_{p}-1<n<n_{p}+1$
Time err. on $\Delta \mathrm{T}$ estimation $[-\mathrm{Tc} \ldots+\mathrm{Tc}] \Rightarrow \Delta \mathrm{T}-\mathrm{Tc}<\mathrm{DTM}<\Delta \mathrm{T}+\mathrm{Tc}$

$$
\mathrm{f}(\Delta \mathrm{~T}-\mathrm{Tc})-1 / \Delta \mathrm{T}<\mathrm{f}_{\mathrm{m}}<\mathrm{f}(\Delta \mathrm{~T}+\mathrm{Tc})+1 / \Delta \mathrm{T}
$$

$$
\mathrm{Ea}_{\text {sup }}=(1+\mathrm{fTc}) / \Delta \mathrm{T}=-\mathrm{Ea}_{\mathrm{inf}} \quad E \mathrm{Er}_{\text {sup }}=\mathrm{Ea} /_{\mathrm{fmin}}=-\mathrm{Er}_{\text {inf }}
$$

## Pulse count during time interval: I-HW, T-int

- $N$ bit counter connected to a microprocessor port, read during interrupt routine at the beginning and at the end of $\Delta \mathrm{T}$. Steps of the interrupt service routine:
- interrupt disable and context saving
- count = present counter value - previous value;
- freq = count/ $\Delta \mathrm{T}$;
- previous value = present counter value;
- context restore and interrupt enable


## -Constraints

$\mathrm{f}_{\text {max }} \Delta \mathrm{T}<2^{\mathrm{N}}$ (f=frequency of pulses); $\mathrm{f}_{\mathrm{P}}=1 / \Delta \mathrm{T} ; \mathrm{f}_{\text {min }}>1 / \Delta \mathrm{T}$

## - Errors

Quantization error on pulses identification [-1 ... 1]
Time error on $\Delta \mathrm{T}$ estimation $\left[-\mathrm{T}_{\mathrm{Imax}} \ldots+\mathrm{T}_{\operatorname{Imax}}\right] \quad \mathrm{T}_{\mathrm{Imax}}=$ max. interrupt latency

$$
\begin{gathered}
\mathrm{f}\left(\Delta \mathrm{~T}-\mathrm{T}_{\operatorname{lmax}}\right)-1 / \Delta \mathrm{T}<\mathrm{f}_{\mathrm{m}}<\mathrm{f}\left(\Delta \mathrm{~T}+\mathrm{T}_{\text {lmax }}\right)+1 / \Delta \mathrm{T} \\
E a_{\text {sup }}=\left(1+\mathrm{f} \mathrm{~T}_{\text {Imax }}\right) / \Delta \mathrm{T}=-\mathrm{Ea}_{\text {inf }} \quad E r_{\text {sup }}=\mathrm{Ea} /_{\text {fmin }}=-\mathrm{Er}_{\text {inf }}
\end{gathered}
$$

## Pulse count during time interval: I-HW, T-HW



## -Constraints

$f_{\max } \Delta \mathrm{T}<2^{\mathrm{N}}$ (f= pulse frequency); $\mathrm{f}_{\mathrm{p}}=1 /\left(\Delta \mathrm{T}+\mathrm{T}_{\text {Imax }}\right) ; \mathrm{f}_{\text {min }}>1 / \Delta \mathrm{T}$

## -Errors

Quantization error on pulses identification [-1 $\ldots$ 1] Time error on $\Delta T$ estimation $\left[-T_{c k} \ldots 0\right]$

$$
\mathrm{f}\left(\Delta \mathrm{~T}-\mathrm{T}_{\mathrm{ck}}\right)-1 / \Delta \mathrm{T}<\mathrm{f}_{\mathrm{m}}<\mathrm{f} \Delta \mathrm{~T}+1 / \Delta \mathrm{T}
$$

## Considerations on pulse count during time interval

-Typical mode: interrupt for time, pulse count with counter

- for a good resolution $\Delta \mathrm{T} \gg$ pulse period
-Constant quantization error on frequency measurement ( $1 / \Delta \mathrm{T}$ )
$\bullet \Delta \mathrm{T}$ has to be chosen considering the minimum frequency to be measured and the relative error


## TRADEOFF ON ERRORS

- $\uparrow \Delta T, \Downarrow$ errors and measurable frequencies $\Downarrow$, $\Uparrow$ measureent delay
$\bullet \downarrow T_{c k}{ }^{\prime} \Downarrow$ errors, $\Uparrow$ timer bits
$\bullet \Downarrow T_{\text {latency int. }} \downarrow$ errors, $\Uparrow$ process management complexity
$\bullet \Downarrow$ Tc $\Downarrow$ errors and measurable frequencies $\Uparrow, \Uparrow$ CPU load


## Measure of the interval between two consecutive pulses

- N bit free-running counter
-Time pulses with duration $\mathbf{T}_{\text {ck }}$ -interrupt to acquire the I and the II pulse -interrupt service routine:

$$
\begin{aligned}
& \text { context_saving } \\
& \text { temp=read counter; } \\
& \text { freq=1/ (temp-prec); } \\
& \text { prec = temp; } \\
& \text { Context_restore }
\end{aligned}
$$

-Constraints

$T_{c k} 2^{N}>1 / f_{\text {min }} ; f_{p}=f$ to measure (minimum $f_{\text {min) }} ; f_{\max }<1 / T_{c k}$

## - Errors

Time quantization error $\left[\begin{array}{lll}-T_{c k} & \ldots & T_{c k}\end{array}\right]$
Time error on pulses identification $\left[-\mathrm{T}_{\operatorname{lmax}} \ldots+\mathrm{T}_{\operatorname{lmax}}\right] \quad \mathrm{T}_{\operatorname{lmax}}=$ max. interrupt latency

$$
\mathrm{f} / 1+\mathrm{f}\left(\mathrm{~T}_{\mathrm{ck}}+\mathrm{T}_{\mathrm{lmax}}\right)<\mathrm{f}_{\mathrm{m}}<\mathrm{f} / 1-\mathrm{f}\left(\mathrm{~T}_{\mathrm{ck}}+\mathrm{T}_{\mathrm{lmax}}\right)
$$

## Considerations on the measurement of the interval between two pulses

- typical mode: that showed
- null frequency measurement is critical
- variable acquisition period ( $f_{p}$ ) and measurement production
- minimum resolution when frequency is high, max. when it is low
- measur. interval $=\mathbf{N}$ pulses, with $\mathbf{N}$ set depending on the chosen resolution
- $\Delta T$ chosen on the basis of the minimum frequency


## COMPROMESSI SU ERRORI

$\Downarrow \mathrm{T}_{\mathrm{ck} \boldsymbol{\prime}} \Downarrow$ errors, $\Uparrow \mathrm{f}_{\text {max }}$ detectable, $\Uparrow$ timer bits
$\Downarrow T_{\text {latency int. }} \Downarrow$ errors, $\Uparrow f_{\text {max }}$ detectable, $\Uparrow$ complexity

## EXAM TEXT (May 2007)

The frequency of a pulse train must be measured in the variable range of [5 $\mathrm{KHz} \ldots 500 \mathrm{KHz}$ ]. The measurement must be provided every $\underline{\mathrm{T}}$ sec, allowing a small tolerance in the periodicity due to the real availability of the CPU:
-Describe how the chosen technique works and choose suitable devices so as the uncertainty range affecting the measure in the worst case is [-0.15\% ... + 0.1\%];
-Verify that the constraints required by the chosen technique are satisfied;
-Determine the maximum and minimum frequencies really measurable with the imposed constraints;
-Calculate the uncertainty range in the frequency measurement if the quantization error on the $\Delta T$ interval estimation is supposed to be equal to +/- "quantum".

## Pulse acquisition for synchronising

- events corresponding to a valid information

- Data_valid; in A, (PORTA_PIO)

- peripheral ready to receive

- Strobe from peripheral; out A, (PORTA_PIO)



## PWM PULSE EMISSION

- Control of processes with a slow behavior that are sensible to the average value of the control signal applied (or of the manipulated variable). For example the environmental temperature
- A pulse with $A$ amplitude and $D(t)$ duration where $0<D(t)<P$ (period)
- Average value $\left(V_{m}\right)=A * D / P$ or $A * D C(t) \quad D C=$ duty cycle (D/P)
- Relative quantization error on $\mathbf{V}_{\mathrm{m}}$ is $\mathrm{P}_{\mathrm{ck}}$ (clock granularity) / $\mathbf{P}$
- We achieve an "analog" signal through a digital one
- If $P$ is in the order of $100 \mathrm{msec}=>$ sw management
- Also to drive dc motors



## EXAM TEXT (July 2006)

- Design a system able to heat up an incubator at $20^{\circ} \mathrm{C}$ for 2 minutes, using a typical microprocessor equipped with a parallel PIO port connected to a buffer and to a amplification stage with gain $A=10$
- Suppose that the ratio between environmental temperature and the average applied power (thermal efficiency) is $80^{\circ} \mathrm{C} / \mathrm{W}$ and that the heater can be electrically modeled as a $1 \mathrm{~K} \Omega$ resistance.
- Suppose to heat for 10 minutes at $30^{\circ}$, for 5 minutes at $20^{\circ}$ and for 5 minutes at $10^{\circ}$ : what is the final temperature reached?



## EXAM TEXT (July 2006)


$W_{E L}=\frac{V^{2}}{R}=0,625 \mathrm{~W}$
Temp $=0,625 \cdot 80=50^{\circ} \mathbf{C}$
$W_{E L_{20^{\circ}} \mathrm{C}}=\frac{20}{80}=0.25 \mathrm{~W}$

$$
V_{20^{\circ} \mathrm{C}}=\sqrt{1000 * 0,25}=\sqrt{250}=15,8
$$

$$
\frac{\Delta T}{120}=\frac{15,8}{25} \rightarrow \Delta T=76,8 s
$$

## EXAM TEXT (July 2006)

$$
\begin{aligned}
& \frac{T 1 \cdot \Delta T_{1}+T 2 \cdot \Delta T_{2}+T 3 \cdot \Delta T_{3}}{\Delta T_{1}+\Delta T_{2}+\Delta T_{3}}=\frac{30 \cdot 10+20 \cdot 5+10 \cdot 5}{20}=22,5^{\circ} \mathrm{C} \\
& T_{1}=30 \rightarrow W_{E L 1}=\frac{30}{80} \rightarrow V_{1}=\sqrt{1000 \cdot \frac{30}{80}}=19,36 \mathrm{~V} \\
& T_{2}=20 \rightarrow W_{E L 1}=\frac{20}{80} \rightarrow V_{2}=15,8 \mathrm{~V} \\
& D_{1,2,3}=\frac{V_{1,2,3}}{V_{\max }} \Delta t_{1,2,3} \\
& T_{3}=10 \rightarrow W_{E L 1}=\frac{10}{80}=0,125 \rightarrow V_{1} \\
& =\sqrt{125}=11,2 \mathrm{~V}
\end{aligned}
$$

## PWM PULSE EMISSION: DC MOTORS

- Transformation of electric energy in mechanic (rotation vel. and torque)
- If the voltage supply is fixed a dc motor absorbs a current proportional to the exerted torque (thus it depends on the applied load)
- When the motor is in stale it absorbs much more current than in the usual case so to avoid damages is important to define two "limit" magnitudes
- Stalemate current: max current absorbed by a motor at the nominal voltage, during the stalemate phase (from hundreds of mA to few A).
- Stalemate torque: the torque provided when the driveshaft is blocked (at the startup or due to an external load), with a nominal voltage and maximum current ( 100 Nm )



## PWM PULSE EMISSION: DC MOTORS

- Sometimes the required velocity is less than that the motor can reach
- Suitable gears can be used (adaptors) that diminish $\omega$ by a suitable constant factor. The ensemble motor + adaptor can be defined as motoadaptor
- The adaptors can be used individually or coupled to regulate the overall angular velocity

$$
\omega_{2} R_{2}=\omega_{1} R_{1}
$$




Single adaptor


$$
\begin{aligned}
& \frac{\vartheta_{1}}{\vartheta_{2}}=\left(\frac{\mathrm{R}_{2}}{\mathrm{R}_{1}}\right)^{2} \\
& \frac{\omega_{1}}{\omega_{2}}=\left(\frac{\mathrm{R}_{2}}{\mathrm{R}_{1}}\right)^{2}
\end{aligned}
$$

Adaptor series

## PWM PULSE EMISSION: DC MOTORS

Let's consider the mechanical power P conservation, $\Gamma \omega=\Gamma^{\prime} \omega^{\prime}$
Thus when the angular velocity drops down due to the effect of an adaptor the exerted torque correspondingly augments

In other words, if $\omega=\mathrm{n} \omega^{\prime}$ then $\Gamma^{\prime}=\mathrm{n} \Gamma$
Finally, since the mechanical behavior of a dc motor can be described by
$\Gamma=\Gamma_{0}+\gamma \omega+J \dot{\omega} \quad \Gamma_{0}$ effective torque, $\gamma$ friction and $J$ inertia
$\Gamma^{\prime}-n \Gamma_{0}=\left(J_{R}+n^{2} J\right) \omega^{\prime}+\left(\gamma_{R}+n^{2} \gamma\right) \omega^{\prime} \quad \Gamma^{\prime}$ new torque, $\gamma_{\mathrm{R}}$ friction and $\mathrm{J}_{\mathrm{R}}$ inertia of the adaptor

Significant increase of inertia and friction Instabilities and hysteretical behavior due to mechanical play The motoadaptors can be effective where high precision is not required

## PWM PULSE EMISSION: DC MOTORS



## Model of the motor without load



## PWM PULSE EMISSION: DC MOTORS

Model of the mechanical part:

$$
\tau=J_{m} \cdot \ddot{\theta}_{m}+B_{m} \cdot \dot{\theta}_{m}+d+\tau_{c}
$$

Similar approach to describe the load components:

$$
\tau^{\prime}=J_{c} \cdot \ddot{\theta}+B_{c} \cdot \dot{\theta}+d^{\prime}
$$

## Model of the motor with load and adaptor

## PWM PULSE EMISSION: DC MOTORS


this brings to $\tau_{c}=J_{c} \cdot \frac{\ddot{\theta}_{m}}{n^{2} \cdot \xi}+B_{c} \cdot \frac{\dot{\theta}_{m}}{n^{2} \cdot \xi}+\frac{d^{\prime}}{n \cdot \xi}$
$\xrightarrow{+}$
The adaptor reduces the effect of disturbs on load
by substituting in the general motor expression

$$
\tau=\left(J_{m}+\frac{J_{c}}{n^{2} \cdot \xi}\right) \cdot \ddot{\theta}_{m}+\left(B_{m}+\frac{B_{c}}{n^{2} \cdot \xi}\right) \cdot \dot{\theta}_{m}+\frac{d^{\prime}}{n \cdot \xi}+d
$$

that is

$$
\tau=J_{e f f} \cdot \ddot{\theta}_{m}+B_{e f f} \cdot \dot{\theta}_{m}-\tau_{d}
$$

where $\quad-\tau_{d}=\frac{d^{\prime}}{n \cdot \xi}+d$

## PWM PULSE EMISSION: DC MOTORS

## Model of the motor with load and adaptor



## PWM PULSE EMISSION: DC MOTORS

## Saturation

(for example due to current limitation)



## Possible non linearities

Adaptor hysteresis (due to possible mechanical plays)


## PWM PULSE EMISSION: DC MOTORS

DC motor, simplified equivalent electric circuit:

- $\mathrm{V}_{\mathrm{a}}=$ supply voltage
- $I_{a}=$ supply current
$\bullet E=$ back electromotive force
- $\mathrm{R}_{\mathrm{a}}=$ armature resistance
- $\mathrm{L}_{\mathrm{a}}=$ armature inductance


Model:

- $V_{a}=E+R_{a} I_{a}$
$\bullet E=k_{e} \omega$
- $P_{a}=V_{a} I_{a}=E I_{a}+R_{a} I_{a}^{2}$

E back emf
$\mathrm{k}_{\mathrm{e}}$ constant for the velocity
absorbed electric power $=$ mechanical power + power lost in the armature

## PWM PULSE EMISSION: DC MOTORS

-DC motor, characteristic:
$P_{m}=$ mechanical power $=E I_{a}=\Gamma_{m} \omega$
$\Gamma_{\mathrm{m}}=$ motor torque $=\mathrm{k}_{\mathrm{t}} \mathrm{I}_{\mathrm{a}}$
-considering that:
$\mathrm{I}_{\mathrm{a}}=\left(\mathrm{V}_{\mathrm{a}}-\mathrm{E}\right) / \mathrm{R}_{\mathrm{a}}$ and $\mathrm{E}=\mathrm{k}_{\mathrm{e}} \omega$ we obtain that $\Gamma_{\mathrm{m}}=\mathrm{k}_{\mathrm{t}}\left(\mathrm{V}_{\mathrm{a}}-\mathrm{k}_{\mathrm{e}} \omega\right) / \mathrm{R}_{\mathrm{a}}$


The torque diminishes when the angular velocity increases
At the beginning the mechanical power grows up with the angular velocity then diminishes with a typical "bell" behavior

## TORQUE-VELOCITY REAL STATIC CHARACTERISTIC




VERY SLOW

## PWM PULSE EMISSION: DC MOTORS

-DC motors, power supply:
Motors absorb a lot of current (0.5:1) A so requiring specific driving circuits (power drivers)

Two possible kinds of driving: linear and on-off


To manage a bi-directional rotation a bipolar power supply $V_{\text {in }}+/$ - is required

If the resistant torque (the ensemble that opposes to rotation inertia, frictions, disturbs, ...) is negligible a good linearity can be obtained between $\mathrm{V}_{\text {in }} \mathrm{e} \omega$

Should not a bipolar supply available, a "bridge" solution can be introduced

## PWM PULSE EMISSION: DC MOTORS

-DC motors, linear driving (a bipolar voltage using a unipolar one):


## PWM PULSE EMISSION: DC MOTORS

-DC motor, on off driving:
If the linear driving dissipates too much power, it can be used only in case of motors that absorb few Watts

To reduce the dissipated power MOS switching devices in on (saturation)/off stage


Dissipated power $=\mathrm{V}_{\mathrm{ce}} \mathrm{I}$
If MOS off, $\mathrm{I}=0$ no dissipated power
If MOS in saturation, $\mathrm{V}_{\mathrm{ce}} \cong 0$ so dissipated power $\cong 0$
"Outflow" diodes for the current are required to avoid excessive voltages at the MOS drain (although it is fast and able to carry high currents)

Limit: the motor will run at the maximum velocity (only mono-direction) otherwise it will be stopped

## PWM PULSE EMISSION: DC MOTORS



## -CC motor, on off driving:

For this reason it is useful to drive the motor through duty cycle variable pulses (PWM) so as to obtain a enough pulse frequency with an average supply voltage proportional to the pulses duty cycle.

However it is still not possible a bi-directional rotation


## PWM PULSE EMISSION: DC MOTORS

-DC motor, H bridge driving (for small laboratory robots):
the " H -bridge" circuit allows to control the sense of rotation.

| Q1 e Q4 on | Clockwise rotation |
| :---: | :---: |
| Q2 e Q3 on | Counterclockwise rot. |
| Q1 e Q2 on | Q1 e Q2 can fire |
| Q1 e Q3 on | Braking |
| All the transistors off | Uncontrolled slowing <br> down |



If two transistors on the same side are active (shoot-through) a big current passes between Vcc and ground. However, now the load is no more the motor (who absorbs the most of the current and so the voltage drop) as before thus the bridge could be damaged.

## PWM PULSE EMISSION: DC MOTORS

-CC motor, H bridge driving:
If we do not want to make the microprocessor serving as pulse emitter, a suitable electronic net could be used to activate transistors
$V_{\text {in }}>V_{d s} Q_{3}$ e $Q_{2}$ active
$V_{\text {in }}<V_{d s} Q_{1}$ e $Q_{4}$ active


Within a period $T$ of the input signal $V_{d s}$ the $V_{d s}$ motor is commanded to rotate for a certain time in a direction and for the remaining in the opposite direction; for the long-term time the resulting direction will be proportional to the sum of the different rotations in the different intervals: if $\mathrm{V}_{\text {in }}$ would be horizontal and exactly located in the medium of $\mathrm{V}_{\mathrm{ds}}$ the motor would be stopped (Locked anti-phase PWM)

## PWM PULSE EMISSION: SERVO MOTORS

They encapsulate dc motor, adaptors, control electronic e position control through a potentiometer for detecting the shaft position and rotation limit switch (to protect the movement area)


Characteristics: torque, rotation velocity, maximum rotation angle. Used in robotics applications


Orders of magnitude:
Rotation angle: $\pm 90^{\circ}$, torque: $3-20 \mathrm{Kg}{ }^{*} \mathrm{~cm}$ (but even 20 Nm ), absorbed current: $1-10$ mA on rest and $\sim 1$ A full load, rotation velocity: up to 9000 rpm , weight: $\sim 150 \mathrm{gr}$.

## PWM PULSE EMISSION: SERVOMOTORS

## -Servomotors:

PCM pulses (Pulse Code Modulation) are used, and the duty cycle determines the shaft position with respect the previous one (the incremental amount of the rotation): for example, depending on the motor type, a certain pulse duration corresponds to $90^{\circ}$ rotation, if the duration is twice the rotation is $+90^{\circ}$, if it is half the rotation is $0^{\circ}$. If the servo does not frequently receive pulses it is abandoned without any control.

Usually the repetition period is 20-30 msec, whilst the pulse duration ranges from a 1 msec minimum up to 2 msec max.
The servo is stopped when a "central" pulse duration is provided or due to its natural lost of energy (frictions ...)
Open loop position control


## PWM PULSE EMISSION: SERVOMOTORS

## -Servomotors:

The "response" of the servo depends mainly on the pulse duration more than on the duty cycle since these motors are designed to properly work into a enough wide pulse frequency range within which the behavior is quite similar (not a orthodox PWM).

If for example a servo receives pulses with variable period among 0-70 msec, it may happen that with values close to $\sim 10 \mathrm{msec}$ the servo becomes a little unstable (vibrations), while when the period is $>40 \mathrm{msec}$, the torque diminishes. Thus a period close to 20 msec could be the right choice (moreover is easy $\Rightarrow 50 \mathrm{~Hz}$ ).
The servo are characterized by a suitable "working curve" to achieve a predetermined motion dynamic.
The "curve" could consist on variable width steps (larger at the beginning, smaller when close to the target) up to reach the set point. It could be useful to control the movement velocity although this is quite difficult since practically the system moves always at maximum velocity.
How to generate the pulse?: a hybrid solution could envisage a $\mu \mathrm{P}$ triggering a timer so as it exits pulses with a duty cycle set again through the $\mu \mathrm{P}$. It is not possible to completely devote a $\mu \mathrm{P}$ to pulse emission because the required times (sometimes < msec ) ask for high frequency routines, with possible $\mu \mathrm{P}$ overload and inaccuracies.

## PWM PULSE EMISSION: DC MOTORS


-Periodic task (15 msec) to set the rotation "step"
-At the beginning the step is wide, when the position is close to the target it diminishes -The position of the rotating shaft can be evaluated time by time

## PFM PULSE EMISSION: STEPPER MOTORS

-Step motors: every received pulse correspond to a precise rotation step
-Low torque, good precision, low cost. Applications: computing systems peripherals, robotics, small applied loads

- Specs:
-Pulse duration enough to drive electronics
-Start frequency (minimum sequence necessary to start rotation) and stop (minimum frequency below which the motor is stopped)
-Acceleration and deceleration ramp
-Possible jitters that imply steps lost or vibrations

-Dedicated hw or processors (for example PIC: programmable interface controller, Microchip Risc microcontroller, low costs)


## PFM PULSE EMISSION: STEPPER MOTORS



How to provide clock?

## PFM PULSE EMISSION: STEPPER MOTORS

| $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | Half step: the step | 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ON | OFF | OFF | OFF | number is twice | ON | OFF | OFF | ON |
| OFF | ON | OFF | OFF | (precision) but the | ON | OFF | OFF | OFF |
| OFF | OFF | ON | OFF | torque is irregular | ON | ON | OFF | OFF |
| OFF | OFF | OFF | ON | and the power | OFF | ON | OFF | OFF |
| ON | OFF | OFF | OFF | consumption not | OFF | ON | ON | OFF |
| Wave: less torque with | constant | OFF | OFF | ON | OFF |  |  |  |
| respect two phase (1vs1.4) |  |  |  |  |  |  |  |  |



Fig 1. Five-wire stepper motor


## PFM PULSE EMISSION: STEPPER MOTORS



Fig 1 - One phase on - full step


Fig3 - One-two phase on - half step


Fig2 - Two phase on - full step


## PFM PULSE EMISSION: STEPPER MOTORS



| 1a. WAVE DRIVE MODE |  | 1b. NORMAL DRIVE MODE |  |
| :---: | :---: | :---: | :---: |
| Transistor Accesi | Posizione Rotore | Transistor Accesi | Posizione Rotore |
| Q1 |  | Q1 e Q2 |  |
| Q2 |  | Q2 e Q3 |  |
| Q3 |  | Q3 e Q4 |  |
| Q4 | $\mathrm{N}(\mathrm{~S} \mathrm{~S} 4$ | Q4 e Q1 |  |

## PFM PULSE EMISSION: STEPPER MOTORS

Interface $\mu \mathbf{P} \mathbf{2 5}$ pin

| X2 pin | Motor | Coil |
| :---: | :---: | :---: |
| 1 | X | A |
| 2 | X | B |
| 3 | X | C |
| 4 | X | D |
| $14,15,16,17$ | X | CENTER-TAP |
| 6 | Y | A |
| 7 | Y | B |
| 8 | Y | C |
| 9 | Y | D |
| $18,19,20,21$ | Y | CENTER-TAP |
| $10,11,12,13$ | INPUTS |  |
| 5 | NOT CONNECTED |  |
| $22,23,24,25$ | POWER SUPPLY GROUND |  |



## PFM PULSE EMISSION: STEPPER MOTORS

## HSTEP motor. <br> .....................Set motor to use half-step drive sequence.

2PHASE motor
Set motor to use two-phase drive sequence. This is the default drive sequence.
WAVE motor. Set motor to wave drive sequence

RATE motor val1 <val2>...Set the delay rate between steps (in milliseconds). Range from 1 to 50,000.
DIR motor <CW, CCW>....Set the direction of rotation. CW = clockwise, CCW = counter-clockwise. If no direction is given then it is reversed.
STEP motor val1 <val2> ....Step motor by val1 steps. If motor is both then val1 refers to ' X " and val2 refers to ' Y '.
WAIT motor.......................Wait for the previous motor command to finish before executing the next command for that motor.
SPIN motor..........................Continuously step motor. Not affected by 'WAIT'.
STOP <motor>
.Immediately stop the current command for motor. If no motor is specified then stops both. Ignores 'WAIT'.
MANUAL ...........................Manually step motors. Another menu appears.
LOOPTIL input(n) HIGH,LOW where $n=1,2,3$ or 4.Wait for the specified input to go high or low before continuing.
DELAY milliseconds..........Waits the specified delay time before continuing. Range = 1 to 50,000.
ECHO ON or OFF .............If ON (default) then all commands are echoed to the output.
PRINT message...................to output
VER......................................Print version number
END....................................Wait for all motor commands to finish then quit
QUIT ....................................Quit program immediately
HELP....................................Display above summary

## STEPPER MOTORS: how to power supply

-This actuator responds with a velocity to the control input variable, allowing to determine the new position without any need of a position transducer (i. e. if we fix a 100 pulses/sec velocity, this corresponds to half a round in a second, if the sensibility of the transducer is 1 round after 200 pulses). Sometimes a limit switch is necessary.
-These motors require power supply in the order of few V , with currents in the order of the $A \Rightarrow r$ in the coils must be relatively low; since the delay is $L / r$ in an inductor, time constants can be high. L moreover can be high if the coils windings are high.

For example if the command voltage is applied to the coil 1 , this one due to the initial 'resistance' will inhibit the, at least at the beginning, the current flow. This means that in the MOS drain $\mathrm{V}_{\mathrm{cc}}$ is nearly present and that the MOS will employ further time to exit from the interdiction towards saturation $\Rightarrow$ this is acceptable if the working frequency is not so high.
It is better to provide higher voltages ( $\Rightarrow \gg$ resistors, < time constants): however this brings to a strong power dispersion on the MOS (heat). Thus a PWM solution is employed with suitable on-off intervals (tunable duty cycle).

## STEPPER MOTORS: PWM POWER SUPPLY



The motor receives a constant current, obtained through a HV voltage necessary to provide fast commutations (time constant $\mathrm{L} / \mathrm{r}$ )
If current is $<\mathrm{V}_{\text {rif }} / \mathrm{R}$ the flip flop close the switch so supplying the motor with a 100 V voltage. Alternatively (open switch), the circuit is grounded through a diode.
The flip flop closes the switch depending on a clock period that allows to regulate the opening and closing intervals (that means the duty cycle).
The inductor serves as "regulation maker" allowing a smooth passage among the levels of the current (high low high low) without shocking the motor.
ON OFF control.

## PULSE EMISSION: WAVE PARTIALIZATION



- "controlled" AC power supply through SCR/TRIAC
- The wave passage through zero must be revealed then a $\mu \mathrm{P}$ should provide the pulses for SCR/TRIAC activation
- Wave period 20 msec , the passage through zero has to be detected with very high precision $\Rightarrow \mathrm{HW}$ solution
- $D(t)$ can be settled through a sw routine or through a timer with a interrupt to $\mu P$
- Non linear relationships between the average power provided and the power up delay


## PULSE EMISSION: WAVE PARTIALIZATION



## PULSE EMISSION: WAVE PARTIALIZATION



$$
\mathrm{T}=\frac{\frac{2}{3} V c c}{I c}
$$

$$
I c=\frac{V c c-V b-0.7}{R}
$$

$$
\mathbf{T}=\frac{2 R C V c c}{3(V c c-V b-0.7)}
$$

$\mathrm{V}_{\mathrm{b}}$ allows to regulate the capacitor charge velocity so to establish when $(\tau)$ to power up the SCR (i. e. the provided power).
High non linearity between the delay and the power applied to the load.

## PULSE EMISSION: WAVE PARTIALIZATION



## CODED SET POINT VALUES



## HW CODED (SETUP) VALUES: BCD SWITCHES

Signals alone do not bring significant information: it's their ordered ensemble that represents an information

A typical example is the acquisition of BCD numeric values for the digital set point of a regulation chain $=$ contraves


Multiswitch contrave with rotational setup of the digits


Multiswitch contrave with linear (push button) setup of the digits


Equivalent electronic circuit of every terminal element (digit)

## HW CODED (SETUP) VALUES: BCD SWITCHES


oem courtesy

Possible problems during commutations:
-Brushing, slithering and contacts rebounds
-Irregular commutations due to long usage
-Passage through halfway not significant values
-Casual order through which the operator moves digits
-Several input instructions (if for example more than 4 BCD digits and $\mu \mathrm{P}$ with $<=16$ bits)

Solutions:
-Ok temporarily "out of range" values
-A suitable "settlement time" (SW) after which values are accepted

- Validation button (HW)

-A sw conversion from the BCD value to the machine binary representation is needed ...


## CODED VALUES: ABSOLUTE ENCODERS

Gray code is used to minimise the error due to a bad alignment of the optics internal to the encoder.

## Pseudocode algorithm

```
commutare := 0
i := N-1
while i >= 0
    BitBin[i] := BitGray[i]
    if commutare = 1
        then BitBin[i] := not BitBin[i]
        commutare := BitBin[i]
    i := i-1
endwhile
```

C version

```
int Gray12_bin (NumGray);
{
    int commuta,i,mask;
        valore = NumGray;
        i = 12;
        mask = 0x800;
        commuta = 0;
        while (i > 0)
    {
        valore ^= commuta;
            commuta = valore & mask;
            commuta >>= 1;
            mask >>= 1;
            i--;
    }
    return (valore);
}
```


## CODED VALUES: ABSOLUTE ENCODERS

## ASSEMBLY VERSION

Routine that takes as input the variable VALORE in Gray code 12 bit and exits with the pure binary converted value in AX (accumulator) register.

```
'GRAY12_BIN:
    MOV CX,12 
CICLO
    XOR [VALORE],BX
    MOV BX,[VALORE]
    AND BX,AX
    SHR AX,1
    SHR BX,1
    LOOP CICLO
    MOV AX,[VALORE]
    RET 13
```

| Integer | Gray code | Binary code |
| :--- | :--- | :--- |
| 0 | 0000 | 0000 |
| 1 | 0001 | 0001 |
| 2 | 0011 | 0010 |
| 3 | 0010 | 0011 |
| 4 | 0110 | 0100 |
| 5 | 0111 | 0101 |
| 6 | 0101 | 0110 |
| 7 | 0100 | 0111 |
| 8 | 1100 | 1000 |
| 9 | 1101 | 1001 |
| 10 | 1111 | 1010 |
| 11 | 1110 | 1011 |
| 12 | 1010 | 1100 |
| 13 | 1011 | 1101 |
| 14 | 1001 | 1110 |
| 15 | 1000 | 1111 |

## OUTPUT CODED VALUES



We suppose that we have to output more bits than those manageable with a single OUT instruction

OUT (PIOA), reg OUT (PIOB), reg1 OUT (PIOC), WR_CMD

## DATA BUS

ADDRESS BUS

